Implementing geometric complexity theory: On the separation of orbit closures via symmetries

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November 10, 2019

Abstract

Understanding the difference between group orbits and their closures is a key difficulty in geometric complexity theory (GCT): While the GCT program is set up to separate certain orbit closures, many beautiful mathematical properties are only known for the group orbits, in particular close relations with symmetry groups and invariant spaces, while the orbit closures seem much more difficult to understand. However, in order to prove lower bounds in algebraic complexity theory, considering group orbits is not enough.

In this paper we tighten the relationship between the orbit of the power sum polynomial and its closure, so that we can separate this orbit closure from the orbit closure of the product of variables by just considering the symmetry groups of both polynomials and their representation theoretic decomposition coefficients. In a natural way our construction yields a multiplicity obstruction that is neither an occurrence obstruction, nor a so-called vanishing ideal occurrence obstruction. All multiplicity obstructions so far have been of one of these two types.

Our paper is the first implementation of the ambitious approach that was originally suggested in the first papers on geometric complexity theory by Mulmuley and Sohoni (SIAM J Comput 2001, 2008): Before our paper, all existence proofs of obstructions only took into account the symmetry group of one of the two polynomials (or tensors) that were to be separated. In our paper the multiplicity obstruction is obtained by comparing the representation theoretic decomposition coefficients of both symmetry groups.

Our proof uses a semi-explicit description of the coordinate ring of the orbit closure of the power sum polynomial in terms of Young tableaux, which enables its comparison to the coordinate ring of the orbit.

Acknowledgements: This paper was written partially when CI was at the Max Planck Institute for Informatics, at the Max Planck Institute for Software Systems, and at the University of Liverpool. This paper contains results that are present in UK's master's thesis at the Universität des Saarlandes.

2012 ACM CCS: Theory of computation \rightarrow Algebraic complexity theory

Keywords: Geometric complexity theory, group orbit, orbit closure, multiplicity obstruction

1 Motivation: Geometric complexity theory

Symmetries The idea of using the symmetries of the determinant $\det_n := \sum_{\pi \in \mathfrak{S}_n} \operatorname{sgn}(\pi) \prod_{i=1}^n x_{i,\pi(i)}$ and the permanent $\operatorname{per}_m := \sum_{\pi \in \mathfrak{S}_m} \prod_{i=1}^m x_{i,\pi(i)}$ to separate algebraic complexity classes was pioneered by Mulmuley and Sohoni in 2001 [MS01]. This approach is

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based on the observation that \det_n and per_m are both characterized by their respective symmetry groups. For example, consider homogeneous degree n polynomials in n^2 variables $x_{1,1}, \ldots, x_{n,n}$. Let X denote the $n \times n$ matrix whose entry in row i and column j is $x_{i,j}$. Then $\det(X) = \det_n$. Now, for matrices $A, B \in \operatorname{SL}_n(\mathbb{C})$ the entries of the matrix AXB are homogeneous linear polynomials in the n^2 variables. The crucial fact is that every homogeneous degree n polynomial q in n^2 variables that satisfies q(AXB) = q(X) equals $\alpha \cdot \det_n$ for some scalar $\alpha \in \mathbb{C}$. This means that \det_n is *characterized by its symmetries*. For the permanent polynomial, an analogous statement holds, and also for many other structurally simpler polynomials, for example for the power sum polynomial $x_1^D + \cdots + x_m^D$ and for the product of variables $x_1 x_2 \cdots x_D$, see [Ike19].

Algebraic complexity theory An affine projection of a polynomial is its evaluation at a point whose coordinates are given by affine linear polynomials, e.g., $(x_1 + x_2 + 1)^2 = x_1^2 + 2x_1x_2 + x_2^2 + 2x_1 + 2x_2 + 1$ is an affine projection of x_1^2 . Kayal proved that it is NP-hard to decide whether a polynomial is an affine projection of another polynomial [Kay12]. Valiant proved [Val79] that every polynomial p is an affine projection of some det_n for n large enough. The smallest n for which this is possible is called the *determinantal complexity* dc(p). The class of sequences of polynomials (p_m) whose sequence of natural numbers $dc(p_m)$ is polynomially bounded is called VP_s. For the permanent we can define the *permanental complexity* pc(p) in a completely analogous manner: pc(p)is the smallest n such that p is an affine projection of per_n. The class of sequences of polynomials (p_m) whose $pc(p_m)$ is polynomially bounded is called VNP. Since $pc(det_n)$ is polynomially bounded, $VP_s \subseteq VNP$. Valiant's flagship conjecture in algebraic complexity theory, which is also known as the *determinant versus permament* conjecture can be succinctly phrased as $VP_s \neq VNP$. This is equivalent to conjecturing that $dc(per_m)$ grows superpolynomially fast.

Homogeneous projections and endomorphism orbits It will be beneficial to phrase Valiant's conjecture in a homogeneous setting: A homogeneous projection of a homogeneous polynomial (i.e., all monomials have the same total degree) is its evaluation at a point whose coordinates are given by homogeneous linear polynomials. The set of all homogeneous projections of det_n to polynomials in the variables x_1, \ldots, x_N can then be written as $\{\det_n(\ell_1, \ldots, \ell_{n^2}) \mid \ell_i \text{ is a homogeneous linear polynomial in } x_1, \ldots, x_N \}$. Note that we put the $n \times n = n^2$ inputs of the determinant in a linear order. The polynomial function $(x_{1,1}, \ldots, x_{n,n}) \mapsto \det_n(\ell_1, \ldots, \ell_{n^2})$ equals the composition $\det_n \circ A$, where A is the linear map $(x_{1,1}, \ldots, x_{n,n}) \mapsto (\ell_1, \ldots, \ell_{n^2})$. As it is common in representation theory, we write $A \cdot \det_n$ or just $A \det_n \text{ or } \det_n \circ A$. The endomorphism orbit $\operatorname{End}_{n^2}\det_n$ is defined as $\{A \det_n \mid A \in \mathbb{C}^{n^2 \times n^2}\}$, which is the set of all homogeneous projections of det_n to polynomials in at most n^2 variables. Since all polynomials in $A \det_n$ are homogeneous of degree n, we have $\operatorname{per}_m \notin \operatorname{End}_{n^2}\det_n$ for any $m \neq n$. This slight technicality is treated by a procedure called padding: For fixed m, n with m < n, define the padded permanent $\operatorname{per}_{m,n} := (x_{n,n})^{n-m} \cdot \operatorname{per}_m$. Let $\operatorname{dc}'(\operatorname{per}_{m,n})$ denote the smallest n such that $\operatorname{per}_{m,n} \in \operatorname{End}_n 2 \det_n$. A short calculation shows that Valiant's conjecture is equivalent to the conjecture that $\operatorname{dc}'(\operatorname{per}_m)$ grows superpolynomially.

Group orbits It turns out that if we restrict $\operatorname{End}_{n^2}\operatorname{det}_n$ to only the points Adet_n for which A is invertible, we get the much simpler group orbit $\operatorname{GL}_{n^2}\operatorname{det}_n := \{g\operatorname{det}_n \mid g \in \operatorname{GL}_{n^2}\} \subseteq \operatorname{End}_{n^2}\operatorname{det}_n$. The group orbit of the determinant consists of "determinants in disguise", i.e., determinants after a base change. The question whether a polynomial p lies in $\operatorname{GL}_{n^2}\operatorname{det}_n$ can be answered in randomized polynomial time [Kay12]. Finding $g \in \operatorname{GL}_{n^2}$ such that $p = g\operatorname{det}_n$ is called the *reconstruction*

problem, which is the focus of several recent papers, where the determinant is replaced with other algebraic computational models, see e.g. [GKQ14], [KS19], [KNS19].

Representation theory Let $H_{\text{det}_n} \subseteq \text{GL}_{n^2}$ denote the symmetry group of det_n. From the viewpoint of algebraic geometry, the set $\text{GL}_{n^2}\text{det}_n$ is an affine variety ([BLMW11, Sec 4.2], [Mat60, Cor., p. 206]) and a homogeneous space that is isomorphic to the quotient $\text{GL}_{n^2}/H_{\text{det}_n}$. It is crucial to note here that the group H_{det_n} does not carry any information about the fact that we study a space of polynomials! In this way, we study the orbit $\text{GL}_{n^2}\text{det}_n$ independently of its embedding into the space of polynomials. This gives a particularly beautiful description of its coordinate ring via invariant theory:

$$\mathbb{C}[\mathrm{GL}_{n^2}\mathrm{det}_n] = \mathbb{C}[\mathrm{GL}_{n^2}]^{H_{\mathrm{det}_n}}$$

where $\mathbb{C}[\operatorname{GL}_{n^2}]^{H_{det_n}}$ is the set of regular H_{\det_n} -invariant functions on the variety GL_{n^2} [BLMW11, (5.2.6)]. The coordinate ring of GL_{n^2} has classically been studied: It is a $\operatorname{GL}_{n^2} \times \operatorname{GL}_{n^2}$ -representation whose representation theoretic decomposition is multiplicity free:

$$\mathbb{C}[\operatorname{GL}_{n^2}] = \bigoplus_{\lambda} \{\lambda\} \otimes \{\lambda^*\},\tag{1.1}$$

where λ runs over all nondecreasing lists of n^2 integers and $\{\lambda\}$ denotes the irreducible rational representation of GL_{n^2} corresponding to λ (see e.g. Section 6 or the textbooks [FH91, Mac95, Ful97] for more information). Eq. (1.1) is known as the algebraic Peter-Weyl theorem. It implies that the multiplicity of λ^* in $\mathbb{C}[\operatorname{GL}_{n^2}\operatorname{det}_n]$ equals the dimension of the H_{det_n} -invariant space $\{\lambda\}^{H_{\operatorname{det}_n}}$. This coefficient $\dim\{\lambda\}^{H_{\operatorname{det}_n}}$ is known as the symmetric rectangular Kronecker coefficient and has been the object of several papers [BOR09, Man11, IP17]. Even though Kronecker coefficients are mysterious, many properties are known (see e.g. [CDW12, Bla12, SS16, IMW17, Liu17, BCMW17] for some recent advances), and character theory is available to compute their values in many cases, see e.g. [Ike12, Appendix]. A better understanding of Kronecker coefficients could lead to a better understanding of $\operatorname{GL}_{n^2}\operatorname{det}_n$, which could eventually help to separate points from $\operatorname{End}_{n^2}\operatorname{det}_n$.

If we replace det_n by other polynomials, we get an analogous theory that is often equally beautiful. The power sum polynomial and the product $x_1 \cdots x_D$ will be of particular interest in this paper. The corresponding coefficients are not Kronecker coefficients, but plethysm coefficients and related coefficients that appear in algebraic combinatorics, see Proposition 4.1 for the power sum and (4.4) for the product of variables.

Closures As we just have seen, group orbits have several desirable properties and can be understood directly via symmetry groups and algebraic combinatorics. But endomorphism orbits do not behave that nicely. In general, $\operatorname{End}_{n^2}\det_n$ is not a variety. In order to enable the study of $\operatorname{End}_{n^2}\det_n$ with methods from algebraic geometry, we go to the closure (Euclidean closure and Zariski closure coincide here by general principles, see [Mum95, §2.C]): $\operatorname{End}_{n^2}\det_n$, which coincides with the group orbit closure $\operatorname{GL}_{n^2}\det_n$, see e.g. [B118, Sec. 3.5]. Hence we have a chain of inclusions $\operatorname{GL}_{n^2}\det_n \subseteq \operatorname{End}_{n^2}\det_n \subseteq \operatorname{GL}_{n^2}\det_n$. The border determinantal complexity $\operatorname{dc}(\operatorname{per}_m)$ is defined as the smallest n such that $\operatorname{per}_{m,n} \in \operatorname{GL}_{n^2}\det_n$. Mulmuley and Sohoni's conjecture (closely related to Bürgisser's conjecture [Bür01, hypothesis (7)]) is that $\operatorname{dc}(\operatorname{per}_m)$ grows superpolynomially. Since $\operatorname{dc}(\operatorname{per}_m) \leq \operatorname{dc}'(\operatorname{per}_m)$, this would imply Valiant's conjecture. Mulmuley and Sohoni's conjecture can be attacked by representation theoretic multiplicities (see Section 3) as follows. If we assume for the sake of contradiction that $\operatorname{GL}_{n^2}\operatorname{per}_{m,n} \subseteq \operatorname{GL}_{n^2}\operatorname{det}_n$, then by Schur's lemma (see [FH91, Lemma 1.7]) the multiplicities satisfy $\operatorname{mult}_{\lambda}\mathbb{C}[\operatorname{GL}_{n^2}\operatorname{per}_{m,n}] \leq \operatorname{mult}_{\lambda}\mathbb{C}[\operatorname{GL}_{n^2}\operatorname{det}_n]$. Thus, if there exists (λ, d)

that satisfies

$$\operatorname{mult}_{\lambda} \mathbb{C}[\overline{\operatorname{GL}_{n^2}\operatorname{per}_{m,n}}]_d > \operatorname{mult}_{\lambda} \mathbb{C}[\overline{\operatorname{GL}_{n^2}\operatorname{det}_n}]_d, \tag{1.2}$$

then $\underline{\mathsf{dc}}(m) > n$. Such a pair (λ, d) is called a *multiplicity obstruction*.

Orbits vs orbit closures The algebraic geometry of $\overline{\operatorname{GL}_{n^2}\operatorname{det}_n}$ and the representation theory of its coordinate ring are rather difficult to understand, see e.g. [Kum15, BHI17]. But the close relationship between orbit and orbit closure gives hope that results can be transferred from the orbit to the closure. Indeed, $\mathbb{C}[\overline{\operatorname{GL}_{n^2}\operatorname{det}_n}] \subseteq \mathbb{C}[\operatorname{GL}_{n^2}\operatorname{det}_n]$ is a subalgebra, and hence we have $\operatorname{mult}_{\lambda}\mathbb{C}[\operatorname{GL}_{n^2}\operatorname{det}_n] \geq \operatorname{mult}_{\lambda}\mathbb{C}[\overline{\operatorname{GL}_{n^2}\operatorname{det}_n}]$. Getting lower bounds on multiplicities in $\mathbb{C}[\overline{\operatorname{GL}_{n^2}\operatorname{det}_n}]$ seems much harder. For example, the result in [Kum15] holds for those *n* for which the Alon-Tarsi property holds, in particular if *n* is an odd prime number ± 1 , see Section 5. The occurrence results in [BIP19] use explicit constructions using Young symmetrizers, which is a laborious process. But as a first step towards lower bounds on $\operatorname{mult}_{\lambda}\mathbb{C}[\overline{\operatorname{GL}_{n^2}\operatorname{det}_n}]$, [BI17] proved that $\operatorname{GL}_{n^2}\operatorname{det}_n$ is open in its closure and that the ring $\mathbb{C}[\operatorname{GL}_{n^2}\operatorname{det}_n]$ is a localization of $\mathbb{C}[\overline{\operatorname{GL}_{n^2}\operatorname{det}_n]$.

2 Our contribution

In this paper we tighten the results from [BI17] in the case of the power sum polynomial. For $m \geq D$ let $p := x_1^D + x_2^D + \cdots + x_m^D$ and let $q := x_1 x_2 \cdots x_D$. Let $G := \operatorname{GL}_m$. For m = D we separate the two families of orbit closures $\overline{Gp} \not\subseteq \overline{Gq}$ of polynomials p and q using multiplicity obstructions λ , i.e., $\operatorname{mult}_{\lambda} \mathbb{C}[\overline{Gp}] > \operatorname{mult}_{\lambda} \mathbb{C}[\overline{Gq}]$. Our key contribution is a proof method that for the first time implements closely the strategy in [MS01, MS08]: Both the lower bound on $\operatorname{mult}_{\lambda} \mathbb{C}[\overline{Gp}]$ and the upper bound on $\operatorname{mult}_{\lambda} \mathbb{C}[\overline{Gq}]$ are obtained directly from the symmetry groups of p and q and the dimension of the spaces of H_{p} - and H_q -invariants in irreducible GL_m -representations. This is the result of our tightening of the relationship between $\operatorname{mult}_{\lambda} \mathbb{C}[\overline{Gp}]$ and $\operatorname{mult}_{\lambda} \mathbb{C}[\overline{Gp}]$, see Theorem 4.3.

Before our paper, all existence proofs of multiplicity obstructions $\overline{Gp} \not\subseteq \overline{Gq}$ for any p and q required to explicitly construct (with multilinear algebra) copies of irreducible representations in $\operatorname{mult}_{\lambda}\mathbb{C}[\overline{Gp}]$. These papers only took into account the symmetry group of q instead of both symmetry groups, see [BI11, BI13b, GIP17, DIP19].

In particular, we prove $\overline{Gp} \not\subseteq \overline{Gq}$ by explicitly constructing a multiplicity obstruction λ in Theorem 4.3 using the symmetry groups of p and q and their representation theoretic decomposition coefficients, but we do *not* construct an explicit function that separates the two orbit closures! Since our obstruction is neither an occurrence obstruction, nor a vanishing ideal occurrence obstruction (see next paragraph), the separating function is quite nontrivial to recover. This is a step in the right direction, since the explicit construction of separating functions for Valiant's conjecture could turn out to be problematic because of the algebraic natural proofs barrier [FSV17, GKSS17, BIJL18].

Occurrence obstructions and vanishing ideal occurrence obstructions

The classical approach of Mulmuley and Sohoni [MS01, MS08] conjectures the existence of so-called occurrence obstructions, which are types λ for which the stronger property $\operatorname{mult}_{\lambda} \mathbb{C}[\overline{\operatorname{Gper}}_{m,n}] > 0 = \operatorname{mult}_{\lambda} \mathbb{C}[\overline{\operatorname{Gdet}}_n]$ holds. These obstructions are not enough to prove strong complexity lower bounds (at least not in the classical setting of det_n vs per_{m,n}), see [IP17, BIP19]. Recently it was shown that there are settings in which multiplicity obstructions are provably stronger than just occurrence obstructions [DIP19]. The types λ that are used in [DIP19] occur in the vanishing ideal of one orbit closure, but not in the vanishing ideal of the other, hence we call them vanishing ideal occurrence obstructions. How useful vanishing ideal occurrence obstructions are for separating orbit closures

is an open question, but it seems unlikely that strong complexity lower bounds can be proved by using only vanishing ideal occurrence obstructions.

The techniques that we develop in this paper study multiplicities and go beyond just occurrence obstructions and vanishing ideal occurrence obstructions. Theorem 4.3 gives the first family of multiplicity obstructions that are neither occurrence obstructions nor vanishing ideal occurrence obstructions, see Proposition 5.3. To prove this fact, we make use of Drisko's and Glynn's progress on the Alon-Tarsi conjecture about the difference between the number of even and odd Latin squares of a given size.

The toy setting as a starting point

Our separation $\overline{Gp} \not\subseteq \overline{Gq}$ with $p = x_1^D + \cdots + x_m^D$ and $q = x_1 \cdots x_D$ is clearly a toy problem, but even though its complexity theoretic relevance is quite limited (multivariate factorization of a power sum) it shares all (as far as we know) crucial geometric and representation theoretic features with the determinant versus permanent problem: Both problems are problems about orbit closures of polynomials, and the group action is the same canonical action. The only difference between the two setups are the specific polynomials p and q. Even though p and q do not share the complexity theoretic properties of the determinant and the permanent, p and q are characterized by their symmetry groups and are stable points (see [BI17]). Therefore this setup can be seen a starting point from which p and q could now be gradually adjusted until some orbit closure separations can be obtained that give lower bounds in algebraic complexity theory.

Structure of the paper

In Section 3 we start with preliminaries that are necessary to state our results precisely in Section 4. The main connection between representation theory and tableau combinatorics is discussed in Sections 6–11. Section 12 proves a technical result about plethysm coefficients that was postponed from Section 4. Section 13 proves the main technical theorem under the assumption that the so-called *Tableau Lifting Theorem* 13.1 is true. The rest of the paper (Sections 14 to 21) is then used to prove the Tableau Lifting Theorem using elementary but subtle Young tableau combinatorics. Even and odd degrees D are treated mostly independently, where the odd degree case is much more involved.

3 Preliminaries

A partition λ is a nonincreasing finite sequence of natural numbers $(\lambda_1, \lambda_2, \ldots)$. We identify a partition with its Young diagram, which is a top-left aligned array of boxes with λ_i boxes in row *i*. For example, the Young diagram for $\lambda = (4, 3, 1)$ is \square . The length $\ell(\lambda)$ is the number of rows in the Young diagram of λ , formally $\ell(\lambda) = \max\{i \mid \lambda_i > 0\}$. The number of boxes of λ is defined as $|\lambda| := \sum_{i \in \mathbb{N}} \lambda_i$. We also define $|\varrho| := \sum_{i \in \mathbb{N}} \varrho_i$ in the case where ϱ is not a partition. The transpose λ^t of a partition λ is the partition corresponding to the Young diagram that is the reflection of λ at the main diagonal, e.g., $(4,3,1)^t = (3,2,2,1)$. The entries of λ^t are the column lengths of λ . For a partition of length $\leq m$ and δ many boxes, we write $\lambda \vdash_m \delta$. For two partitions λ and μ we define their sum $\lambda + \mu$ in a row-wise fashion: $(\lambda + \mu)_i := \lambda_i + \mu_i$. For natural numbers a, b let $a \times b$ denote the partition that corresponds to the rectangular Young diagram with a rows and b columns, i.e., $a \times b := (b, b, \ldots, b)$.

Fix $m \in \mathbb{N}$. For a partition $\lambda \vdash_m \delta$ we write $\{\lambda\}$ to denote the irreducible GL_m -representation of type λ . These $\{\lambda\}$ form a pairwise non-isomorphic list of irreducible polynomial GL_m representations, see [Ful97]. The dual (=contragredient) representation of $\{\lambda\}$ is denoted by $\{\lambda^*\}$. Since G is a reductive group, every finite dimensional G-representation \mathscr{V} can be decomposed into a direct sum of irreducible G-representations, and we write $\operatorname{mult}_{\lambda}\mathscr{V}$ for the multiplicity of $\{\lambda\}$ in such a decomposition (although the decomposition might not be unique, the multiplicity is the same in any decomposition).

A tableau of shape λ over some finite alphabet \mathcal{A} is a mapping $\lambda \to \mathcal{A}$ of boxes to elements in \mathcal{A} . For example a tableau of shape (4,3,1) over the alphabet \mathbb{N} is given by $\frac{|\frac{1}{3}|\frac{1}{2}|2|}{|\frac{1}{3}|\frac{1}{2}|}$. For a tableau T let $\mathsf{sh}(T) := \lambda$ denote its *shape*, i.e., its vector of row lengths. For a tableau T of shape $\lambda \vdash_m \delta$ over the alphabet $\{1, \ldots, m\}$ we define its *content* to be the vector $\varrho \in (\mathbb{N}_{\geq 0})^m$ where ϱ_i counts the number of occurrences of i in T. For example, the tableau $\frac{1}{2}\frac{3}{2}\frac{3}{2}\frac{1}{2}}{2}$ has content (2, 4, 2). The *sorted content* of a tableau is the partition obtained by sorting the entries in the content in a decreasing order, e.g. (4, 2, 2) in the preceding example. For two tableaus T and S we define their *concatenation* T + S by concatenating rows. The resulting tableau T + S has shape $\mathsf{sh}(T) + \mathsf{sh}(S)$. For example, $\frac{1}{2}\frac{1}{2}\frac{1}{2}\frac{1}{2}\frac{1}{2}} + \frac{4}{5}\frac{4}{5}\frac{1}{6}} = \frac{1}{2}\frac{1}{2}\frac{1}{2}\frac{1}{2}\frac{1}{5}\frac{1}{5}}$. A tableau with entries from \mathbb{N} is called *semistandard* if the entries of each column are strictly increasing from top to bottom. A semistandard tableau T is called *regular* if it does not have a repeated entry. A tableau T is called *regular* if each of its columns is regular. A tableau L is called *duplex* if each column in L appears an even number of times. For example, $L = \frac{1}{2}\frac{1}{2}\frac{1}{2}\frac{1}{3}\frac{1}{3}\frac{1}{4}\frac{1}{4}$ is duplex, while $L = \frac{1}{2}\frac{1}{2}\frac{1}{2}\frac{1}{3}\frac{1}{3}\frac{1}{4}\frac{1}{4}$ is not duplex. Duplex tableaux are the main idea behind the construction in [BCI11], see also [Ike12, Sec. 6.2].

Let $\operatorname{Poly}^{D}\mathbb{C}^{m}$ denote the vector space of homogeneous degree D polynomials in m variables. Let $G := \operatorname{GL}_{m}$. The group G acts on $\operatorname{Poly}^{D}\mathbb{C}^{m}$ via $(gp)(x) := p(g^{t}x)$ for all $g \in G$, $p \in \operatorname{Poly}^{D}\mathbb{C}^{m}$, $x \in \mathbb{C}^{m}$, where the transpose makes this a left action. We consider the power sum polynomial $p := x_{1}^{D} + \cdots + x_{m}^{D} \in \operatorname{Poly}^{D}\mathbb{C}^{m}$. Clearly p is fixed by permuting variables and by rescaling any variable by Dth roots of unity. For $D \geq 3$ these group elements generate the whole stabilizer of p, see e.g. [CKW10]. Let $H := \operatorname{stab} p = \mathbb{Z}_{D}^{m} \rtimes \mathfrak{S}_{m} \subseteq G$ denote the stabilizer of p. Let $q := x_{1} \cdots x_{D}$. Clearly q is fixed under permuting the variables and under rescaling each variable by a scalar α_{i} such that their product $\prod_{i=1}^{D} \alpha_{i}$ equals 1. These elements generate the stabilizer of q, see e.g. [Ike19, Prop. 3.1].

Considering $\operatorname{Poly}^D \mathbb{C}^m$ as a vector space with a *G*-action, let $\mathbb{C}[\operatorname{Poly}^D \mathbb{C}^m]_d$ be the vector space of homogeneous degree *d* polynomials on $\operatorname{Poly}^D \mathbb{C}^m$. In a natural way, $\mathbb{C}[\operatorname{Poly}^D \mathbb{C}^m]_d$ is a *G*representation: $(gf)(q) := f(g^{-1}q)$ for all $g \in G$, $f \in \mathbb{C}[\operatorname{Poly}^D \mathbb{C}^m]_d$, $q \in \operatorname{Poly}^D \mathbb{C}^m$. The multiplicity $\operatorname{mult}_{\nu^*} \mathbb{C}[\operatorname{Poly}^D \mathbb{C}^m]_d$ is called the *plethysm coefficient*, denoted by $a_{\nu}(d, D)$. Note that it does not depend on *m* for $m \geq \ell(\lambda)$, see e.g. [Ike12, Sec. 4.3]. For the empty partition (0) we define $a_{(0)}(0, i) = 1$. Whether or not the plethysm coefficient has a nice combinatorial description is an open research question in algebraic combinatorics, see Problem 9 in [Sta00]. Among computer scientists, this question is commonly phrased as whether or not the map $(\nu, d, D) \mapsto a_{\nu}(d, D)$ is in the complexity class $\#\mathsf{P}$.

Given two irreducible G-representations $\{\mu\}$ and $\{\nu\}$, their tensor product $\{\mu\} \otimes \{\nu\}$ is a $G \times G$ -representation. Embedding $G \hookrightarrow G \times G$ diagonally via $g \mapsto (g, g)$ the tensor product $\{\mu\} \otimes \{\nu\}$ becomes a G-representation that decomposes into irreducibles. The multiplicity $\operatorname{mult}_{\lambda}\{\mu\} \otimes \{\nu\}$ is called the *Littlewood-Richardson coefficient* $c_{\mu,\nu}^{\lambda}$. This quantity has numerous beautiful combinatorial interpretations, see e.g. [BZ92], [Ful97, Sec. 5], [KT99], [Buc00], [Ike12, Sec. 10] and many more. In particular, the map $(\lambda, \mu, \nu) \mapsto c_{\mu,\nu}^{\lambda}$ is in the complexity class $\#\mathsf{P}$. Even though the exact computation of $c_{\mu,\nu}^{\lambda}$ is NP-hard [Nar06], deciding its positivity is possible in polynomial time

[DLM06], [MNS12], [BI13a]. Completely analogous properties hold when we take tensor products of polynomially many irreducible *G*-representations $\{\mu^1\} \otimes \cdots \otimes \{\mu^d\}$ and embed $G \hookrightarrow G \times G \times \cdots \times G$. The corresponding coefficient is called the *multi-Littlewood-Richardson coefficient* $c^{\lambda}_{\mu^1,\mu^2,\ldots,\mu^d}$.

For a partition $\rho \vdash_m d$ the frequency notation $\hat{\rho} \in \mathbb{N}^m$ is defined via $\hat{\rho}_i := |\{j \mid \rho_j = i\}|$. For example, the frequency notation of $\rho = (3, 3, 2, 0)$ is $\hat{\rho} = (0, 1, 2, 0)$. We observe that $|\rho| = \sum_i i \hat{\rho}_i$.

4 Result details

The following Proposition 4.1 writes the multiplicity $\operatorname{mult}_{\lambda^*} \mathbb{C}[Gp]$ as a nonnegative sum of products of multi-Littlewood-Richardson coefficients and plethysm coefficients.

4.1 Proposition. Let $\lambda \vdash_m dD$. Define

$$b(\lambda, \varrho, D, d) := \sum_{\substack{\mu^1, \mu^2, \dots, \mu^d \\ \mu^i \vdash Di\hat{\varrho}_i}} c^{\lambda}_{\mu^1, \mu^2, \dots, \mu^d} \prod_{i=1}^d a_{\mu^i}(\hat{\varrho}_i, iD).$$

Then

$$\mathrm{mult}_{\lambda^*}\mathbb{C}[Gp] = \sum_{\varrho \vdash_m d} b(\lambda, \varrho, D, d)$$

The proof technique is based on the technique in [BI11]. The proof is postponed to Section 10.

We remark that if Problem 9 in [Sta00] is resolved positively, then Proposition 4.1 implies that the multiplicity $\operatorname{mult}_{\lambda^*} \mathbb{C}[Gp]$ has a combinatorial description, i.e., the map $(\lambda, m, d, D) \mapsto \operatorname{mult}_{\lambda^*} \mathbb{C}[Gp]$ is in $\#\mathbb{P}$. The same holds also for its summands $b(\lambda, \varrho, D, d)$. It is known that $\operatorname{mult}_{\lambda^*} \mathbb{C}[Gq] = a_{\lambda}(D, d)$ (see e.g. [Lan17, Sec. 9.2.3]), so the same holds for $\operatorname{mult}_{\lambda^*} \mathbb{C}[Gq]$.

Our main technical theorem that enables us to find obstructions is the following.

4.2 Theorem (Main technical theorem). Let $m, d, D \in \mathbb{N}$. If D is odd, we assume that $\binom{2(D-1)}{D-1} \geq 2(m-1)$. Let $\lambda \vdash_m dD$. For each $\varrho \vdash_m d$ define the number e_{ϱ} as follows:

- if D is even, let $e_{\varrho} := \sum_{i=1}^{m} \left\lceil \frac{\varrho_i}{D-2} \right\rceil$.
- if D is odd, let $e_{\varrho} := \sum_{i=1}^{m} 2 \lceil \frac{\varrho_i}{2(D-2)} \rceil$.

Let Ξ be a subset of the set of all partitions $\rho \vdash_m d$. Let $e_{\Xi} := \max\{e_{\rho} \mid \rho \in \Xi\}$. Then

$$\mathrm{mult}_{(\lambda + (m \times e_{\Xi}D))^*} \mathbb{C}[\overline{Gp}] \geq \sum_{\varrho \in \Xi} b(\lambda, \varrho, D, d).$$

The proof of Theorem 4.2 is postponed to Section 13.

Explicit obstructions via symmetries

We use Theorem 4.2 as follows to construct obstructions.

4.3 Theorem. Let $m = D \ge 4$ be even. Let d = 2. Let $\lambda = (2m)$ and let $\nu = (2m) + (m \times 2m)$. Let $p := x_1^m + \cdots + x_m^m$ and $q := x_1 x_2 \cdots x_m$. Let $\Xi = \{(2), (1, 1)\}$. In the notation of Theorem 4.2 we have $e_{(2)} = 1$ and $e_{(1,1)} = 2$, thus $e_{\Xi} = 2$. Note that $\nu = \lambda + (m \times e_{\Xi}D)$. We have

• b((2m), (2), m, 2) = 1 and b((2m), (1, 1), m, 2) = 1 and hence with Theorem 4.2 we have $\operatorname{mult}_{\nu^*} \mathbb{C}[\overline{Gp}] \ge 2.$

• $1 \geq \operatorname{mult}_{\nu^*} \mathbb{C}[Gq] \geq \operatorname{mult}_{\nu^*} \mathbb{C}[\overline{Gq}].$

In particular $\operatorname{mult}_{\nu^*}\mathbb{C}[\overline{Gp}] \geq 2 > 1 \geq \operatorname{mult}_{\nu^*}\mathbb{C}[\overline{Gq}]$ and hence ν is a multiplicity obstruction that proves the separation $\overline{Gp} \not\subseteq \overline{Gq}$.

Proof. We use the following well-known properties about Littlewood-Richardson coefficients (see e.g. [Ful97, Ch. 5]).

• The Littlewood-Richardson coefficient is symmetric in the subscript parameters: $c_{\mu,\nu}^{\lambda} = c_{\nu,\mu}^{\lambda}$.

• $c_{\mu,(0)}^{\lambda}$ equals 1 iff $\mu = \lambda$, otherwise $c_{\mu,(0)}^{\lambda} = 0$. Recall that $\hat{\varrho}$ is the frequency notation of ϱ . We calculate

$$\begin{split} b(\lambda,(2),m,2) &= \sum_{\substack{\mu^1,\mu^2\\\mu^i\vdash mi(\widehat{2})_i}} c_{\mu^1,\mu^2}^{\lambda} \prod_{i=1}^2 a_{\mu^i}(\widehat{(2)}_i,im) \\ &= \sum_{\mu^2\vdash 2D} c_{(0),\mu^2}^{\lambda} \cdot a_{(0)}(\widehat{(2)}_1,m) \cdot a_{\mu^2}(\widehat{(2)}_2,2m) \underbrace{=}_{c_{(0),\mu}^{\lambda} = \delta_{\mu,\lambda}} a_{(0)}(\widehat{(2)}_1,m) \cdot a_{\lambda}(\widehat{(2)}_2,2m) \\ &= a_{(0)}(0,m) \cdot a_{(2m)}(1,2m) = 1 \cdot 1 = 1. \end{split}$$

Here we used the trivial fact that $a_{(\delta)}(1,\delta) = 1$ for all δ .

$$\begin{split} b(\lambda,(1,1),m,2) &= \sum_{\substack{\mu^{1},\mu^{2} \\ \mu^{i}\vdash mi(\widehat{1,1})_{i}}} c_{\mu^{1},\mu^{2}}^{\lambda} \prod_{i=1}^{2} a_{\mu^{i}}(\widehat{(1,1)}_{i},im) \\ &= \sum_{\mu^{1}\vdash m} c_{\mu^{1},(0)}^{\lambda} \cdot a_{\mu^{1}}(\widehat{(1,1)}_{1},m) \cdot a_{(0)}(\widehat{(1,1)}_{2},2m) \underbrace{=}_{c_{\mu,(0)}^{\lambda} = \delta_{\mu,\lambda}} a_{\lambda}(\widehat{(1,1)}_{1},m) \cdot a_{(0)}(\widehat{(1,1)}_{2},2m) \\ &= a_{(2m)}(2,m) \cdot a_{(0)}(0,2m) = 1 \cdot 1 = 1. \end{split}$$

Here we used the classical fact that $a_{\lambda}(2,m) = 1$ if λ has 2m boxes and at most 2 rows and both rows have even length (see the formula for $h_2 \circ h_n$ in [Mac95, I.8, Exa. 9(a), p. 140]). This proves the first part of the claim.

 $\mathbb{C}[\overline{Gq}] \subseteq \mathbb{C}[Gq]$ is a subalgebra, so $\operatorname{mult}_{\nu^*}\mathbb{C}[\overline{Gq}] \leq \operatorname{mult}_{\nu^*}\mathbb{C}[Gq]$. To show $\operatorname{mult}_{\nu^*}\mathbb{C}[Gq] \leq 1$ we use that

$$\operatorname{mult}_{\nu^*} \mathbb{C}[Gq] = a_{\nu}(m, 2m+2), \tag{4.4}$$

see e.g. [Lan17, Sec. 9.2.3]. We apply the upper bound given by the Kostka numbers $a_{\nu}(m, 2m +$ $2 \leq K(\nu, m \times (2m+2))$, which is the number of semistandard tableaux of shape ν and content $m \times (2m+2)$. This classical upper bound can be deduced for example from [Gay76], see also [Ike12, Sec. 4.3(A)]. It is clear from the special shape of ν that this Kostka number is 1. As an illustration, we give an example of this tableau in the case where m = 4:



In Section 5 we prove that the obstructions in Theorem 4.3 are neither occurrence obstructions nor vanishing ideal occurrence obstructions (in infinitely many cases. This holds in all cases if the Alon-Tarsi conjecture is true).

A remark on plethysm coefficients As far as we know, Theorem 4.2 is the first result of its type in the literature so far. Even the following direct corollary about plethysm coefficient positivity is new.

4.5 Corollary. Let $D \ge 3$ be odd and let m be arbitrary with $\binom{2(D-1)}{D-1} \ge 2(m-1)$. Let d be arbitrary. Let $e = 2\lceil \frac{d}{2(D-2)} \rceil$. Then $a_{(dD)+m \times eD}(d+me, D) \ge 1$.

Proof. Let $\lambda = (dD)$ and $\varrho = (d)$, $\Xi = \{(d)\}$. Theorem 4.2 implies $\operatorname{mult}_{((dD)+(m \times eD))^*} \mathbb{C}[\overline{Gp}] \ge b((dD), (d), D, d) = 1$. Since \overline{Gp} is a subvariety of $\operatorname{Poly}^D \mathbb{C}^m$, it follows $a_{(dD)+m \times eD}(d+me, D) \ge 1$.

Many additional direct corollaries of this type can be drawn from Theorem 4.2.

5 Neither occurrence obstructions nor vanishing ideal occurrence obstructions

The main new property of our obstructions in Theorem 4.3 is that they use both the symmetry group of p and the symmetry group of q. This is a fundamentally new way of constructing obstructions. To highlight the novelty of the approach, in this section we prove that the obstructions in Theorem 4.3 are not vanishing ideal occurrence obstructions. We also prove that they are not occurrence obstructions, provided a property of Latin squares is true (which we know is true for an infinite number of cases). The novelty of our method gives hope that more and stronger results can be proved in a similar way.

Recall the following proposition from [BI17]:

5.1 Lemma ([BI17]). If D is even, then $\operatorname{mult}_{\lambda^*} \mathbb{C}[Gp] = \operatorname{mult}_{(\lambda + (m \times D))^*} \mathbb{C}[Gp]$. If D is odd, then $\operatorname{mult}_{\lambda^*} \mathbb{C}[Gp] = \operatorname{mult}_{(\lambda + (m \times 2D))^*} \mathbb{C}[Gp]$.

5.2 Lemma. Let $\lambda \vdash_m dD$. If D is even, let

$$e := \begin{cases} d & \text{if } d \le m \\ m + \lfloor \frac{d-m}{D-2} \rfloor & \text{if } d \ge m \end{cases}$$

If D is odd and $\binom{2(D-1)}{D-1} \ge 2(m-1)$, let

$$e := \begin{cases} 2d & \text{if } d \le m \\ 2m + 2\lfloor \frac{d-m}{2(D-2)} \rfloor & \text{if } d \ge m \end{cases}.$$

In both cases we have $\operatorname{mult}_{(\lambda+(m\times eD))^*}\mathbb{C}[\overline{Gp}] = \operatorname{mult}_{\lambda^*}\mathbb{C}[Gp] = \operatorname{mult}_{(\lambda+(m\times eD))^*}\mathbb{C}[Gp].$

Proof. The second equality follows from Lemma 5.1.

Let Ξ denote the set of all partitions $\rho \vdash_m d$. Using Theorem 4.2, then according to Proposition 4.1 we have

$$\operatorname{mult}_{(\lambda + (m \times e_{\Xi}D))^*} \mathbb{C}[\overline{Gp}] \ge \operatorname{mult}_{\lambda^*} \mathbb{C}[Gp].$$

Using Lemma 5.1 (and the fact that e_{Ξ} is even if D is odd) we see that

$$\operatorname{mult}_{(\lambda + (m \times e_{\Xi}D))^*} \mathbb{C}[\overline{Gp}] = \operatorname{mult}_{\lambda^*} \mathbb{C}[Gp]$$

It remains to show that $e_{\Xi} = e$. Recall that for natural numbers a, b we have $\left\lfloor \frac{1+a}{b} \right\rfloor - 1 = \lfloor \frac{a}{b} \rfloor$.

Let *D* be even and $d \leq m$. Then the number of nonzero ρ_i is at most *d*. Hence $e_{\rho} \leq d$. On the other hand, $\rho = (1, 1, \ldots, 1, 0, 0, \ldots, 0) \vdash_m d$ provides $e_{\rho} = d$, so the bound is tight. The argument for *D* odd and $d \leq m$ is completely analogous and yields e = 2d.

Let D be even and $d \ge m$. Then $\varrho = (1 + d - m, 1, 1, \dots, 1) \vdash_m d$ provides $e_{\varrho} = m + (\lceil \frac{1 + d - m}{D - 2} \rceil - 1) = m + \lfloor \frac{d - m}{D - 2} \rfloor$. The upper bound is provided via

$$e_{\varrho} = \sum_{i=1}^{m} \left\lceil \frac{\varrho_i}{D-2} \right\rceil = \sum_{i=1}^{m} \left(\left\lceil \frac{\varrho_i + 1 - 1}{D-2} \right\rceil + 1 - 1 \right) = \sum_{i=1}^{m} \left(\left\lfloor \frac{\varrho_i - 1}{D-2} \right\rfloor + 1 \right) = m + \underbrace{\sum_{i=1}^{m} \left\lfloor \frac{\varrho_i - 1}{D-2} \right\rfloor}_{\leq \lfloor \frac{d-m}{D-2} \rfloor} \leq m + \lfloor \frac{d-m}{D-2} \rfloor.$$

Let D be odd and $d \ge m$. Then $\varrho = (1+d-m, 1, 1, \dots, 1) \vdash_m d$ provides $e_{\varrho} = 2m + 2(\lceil \frac{1+d-m}{2(D-2)} \rceil - 1) = 2m + 2\lfloor \frac{d-m}{2(D-2)} \rfloor$. The upper bound is provided via

$$e_{\varrho} = \sum_{i=1}^{m} 2 \lceil \frac{\varrho_i}{2(D-2)} \rceil = \sum_{i=1}^{m} 2 \left(\lceil \frac{\varrho_i + 1 - 1}{2(D-2)} \rceil + 1 - 1 \right) = \sum_{i=1}^{m} 2 \left(\lfloor \frac{\varrho_i - 1}{2(D-2)} \rfloor + 1 \right)$$

= $2m + 2 \underbrace{\sum_{i=1}^{m} \lfloor \frac{\varrho_i - 1}{2(D-2)} \rfloor}_{\leq \lfloor \frac{d-m}{2(D-2)} \rfloor} \leq 2m + 2 \lfloor \frac{d-m}{2(D-2)} \rfloor.$

A Latin square of dimension m is an $m \times m$ matrix for which in each row and in each column each number $1, \ldots, m$ appears exactly once. The *sign* of a column is 1 if the permutation in the column is even, and -1 otherwise. The *sign* of a Latin square is defined as the product of all column signs. A Latin square is called *even* if its sign is 1, and *odd* otherwise. If m is odd, then it is easy to construct an involution on the set of all $m \times m$ Latin squares that pairs each even Latin square with an odd one. Hence, if m is odd, then the number of even $m \times m$ Latin squares equals the number of odd $m \times m$ Latin squares. If m is even, then Alon and Tarsi [AT92] conjecture that the number of even $m \times m$ Latin squares differs from the number of odd $m \times m$ Latin squares (see also [HR94] and [KL15] for equivalent formulations). This is proved for all $m = \tau + 1$ [Dri97] for an odd prime number τ and for all $m = \tau - 1$ [Gly10] for an odd prime number τ , making m = 26the smallest open case. If the number of even $m \times m$ Latin squares differs from the number of odd $m \times m$ Latin squares, then we say that m satisfies the Alon-Tarsi condition. The Alon-Tarsi condition first appeared in connection with geometric complexity theory in [Kum15].

5.3 Proposition. Let ν , m, D, d be as in Theorem 4.3.

- $\operatorname{mult}_{\nu^*}\mathbb{C}[\overline{Gp}] < a_{\nu}(2(m+2),m)$ and hence ν is not a vanishing ideal occurrence obstruction.
- If m satisfies the Alon-Tarsi condition, then $\operatorname{mult}_{\nu^*} \mathbb{C}[\overline{Gq}] > 0$ and hence ν is not an occurrence obstruction. In particular this is true for $m = \tau \pm 1$ for all odd primes τ .

Proof. The first bullet point is treated as follows. $\operatorname{mult}_{\nu^*}\mathbb{C}[\overline{Gp}] = \operatorname{mult}_{\nu^*}\mathbb{C}[Gp]$ by Lemma 5.2. Moreover, $\operatorname{mult}_{\nu^*}\mathbb{C}[Gp] = 2$ (Proposition 4.1 and Theorem 4.3). It remains to prove that $a_{\nu}(2(m+1), m) \geq 3$, which is postponed to Proposition 12.1.

The second bullet point is treated as follows. Kumar proved [Kum15] that $\operatorname{mult}_{(m)^*} \mathbb{C}[\overline{G(x_1 \cdots x_m)}] \geq 1$ and that $\operatorname{mult}_{(m \times m)^*} \mathbb{C}[\overline{G(x_1 \cdots x_m)}] \geq 1$, provided that m satisfies the Alon-Tarsi condition. Since $\nu = (m \times m) + (m \times m) + (m) + (m)$, the semigroup property for occurrences in $\mathbb{C}[\overline{G(x_1 \cdots x_m)}]$ (see e.g. [DIP19]) yields that $\operatorname{mult}_{\nu^*} \mathbb{C}[\overline{G(x_1 \cdots x_m)}] \geq 1$.

The rest of this paper (besides Section 12) is dedicated to the proof of Theorem 4.2.

6 Preliminaries - Representation Theory

In the remainder of this paper we write $G := \operatorname{GL}_m(\mathbb{C})$ to denote the general linear group for some fixed natural number m. A representation of G is a finite dimensional complex vector space \mathscr{V} together with a group homomorphism $\xi : G \to \operatorname{GL}(\mathscr{V})$. If ξ is given by a polynomial map, then we call (\mathscr{V}, ξ) a polynomial representation. We write $gf := \xi(g)(f)$ for all $g \in G$ and $f \in \mathscr{V}$. A linear subspace $\mathscr{W} \subseteq \mathscr{V}$ that is closed under the action of G is called a subrepresentation. A representation \mathscr{V} is called *irreducible* iff it has only two subrepresentations, namely the zero-dimensional subspace and \mathscr{V} itself.

For a G-representation \mathscr{V} , a highest weight vector $f \in \mathscr{V}$ of weight $\lambda \in \mathbb{Z}^m$ is defined to be a vector that satisfies the following two properties:

- for all upper triangular matrices g with 1s on the main diagonal we have gf = f, and
- for all diagonal matrices $g := \text{diag}(a_1, \ldots, a_m)$ we have $gf = a_1^{\lambda_1} \ldots a_m^{\lambda_m} f$.

We denote by $\text{HWV}_{\lambda}(\mathscr{V})$ the vector space of highest weight vectors of weight λ in \mathscr{V} . It turns out that for an irreducible polynomial *G*-representation \mathscr{V} there is exactly one partition $\lambda \in \mathbb{N}^m$ such that $\text{HWV}_{\lambda}(\mathscr{V})$ has dimension = 1, while for all $\mu \neq \lambda$ we have dim $\text{HWV}_{\mu}(\mathscr{V}) = 0$. In this case we write $\mathscr{V} = \{\lambda\}$. Moreover, for each partition $\lambda \vdash_m$ there is an irreducible polynomial *G*-representation $\{\lambda\}$ and we call it the irreducible *G*-representation of *isomorphism type* λ . Furthermore, these $\{\lambda\}$ are pairwise non-isomorphic. We can now define the *multiplicity* $\text{mult}_{\lambda}(\mathscr{V})$ of $\{\lambda\}$ in \mathscr{V} as

$$\operatorname{mult}_{\lambda}(\mathscr{V}) = \operatorname{dim} \operatorname{HWV}_{\lambda}(\mathscr{V}),$$

which is the same as the multiplicity with which $\{\lambda\}$ appears as a summand in a decomposition of \mathscr{V} into a direct sum of irreducible *G*-representations.

The irreducible *G*-representations

In the following exposition we closely follow [Ful97].

For a tablean T let T(r, c) denote the entry of T in row r and column c. Let $(e_i)_i$ be the standard basis of \mathbb{C}^m . Let $\mu = \lambda^t$.

To each tableau $T:\lambda\to\{1,\ldots,m\}$ we assign a tensor

$$e_{T(1,1)} \otimes e_{T(2,1)} \otimes \cdots \otimes e_{T(\mu_1,1)} \otimes e_{T(1,2)} \otimes e_{T(2,2)} \otimes \cdots \otimes e_{T(\mu_2,2)} \otimes \cdots \otimes e_{T(\mu_{\lambda_1},\lambda_1)}.$$

These form a basis of $\bigotimes^{|\lambda|} \mathbb{C}^m$. In this way, the space $\bigotimes^{|\lambda|} \mathbb{C}^m$ is isomorphic to the space of formal linear combinations of tableaux $T : \lambda \to \{1, \ldots, m\}$. Using this isomorphism, the space of tableaux inherits the natural *G*-action on $\bigotimes^{|\lambda|} \mathbb{C}^m$, which is given by

$$g(v_1 \otimes \cdots \otimes v_{|\lambda|}) = g(v_1) \otimes \cdots \otimes g(v_{|\lambda|}).$$

The following vectors span the linear subspace $K(\lambda)$ of *Grassmann-Plücker relations* (sometimes called *shuffle relations*), which is invariant under the *G*-action:

- T + T', where T' is a tableau that arises from T by switching two numbers within one column.
- $T \Sigma S$, where for two fixed columns j, j' and a fixed number of entries k the sum is over all tableaux S that arise from T by exchanging the top k entries in column j with any k entries in column j', preserving the internal vertical order.

It turns out that $\{\lambda\} \simeq (\bigotimes^{|\lambda|} \mathbb{C}^m)/K(\lambda)$, see [Ful97, Ch. 8]. A basis of $\{\lambda\}$ is given by the semistandard tableaux of shape λ with entries from $\{1, \ldots, m\}$. The unique highest weight vector (up to scale) in $\{\lambda\}$ is the *superstandard tableau* of shape λ , which is the semistandard tableau that has only entries *i* in row *i*. It has weight λ .

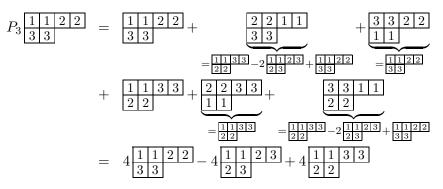
7 Two similar projections of tableaux

In this section we present two symmetrizations that look quite similar on the basis of tableaux. We crucially use this similarity in Section 13 in the proof of the Main Theorem 4.2. Indeed, this peculiarity is the driving force behind our result.

We embed $\mathfrak{S}_m \subseteq G$ via permutation matrices. Given a tableau $S : \lambda \to \{1, \ldots, m\}$ we define:

$$P_m S := \sum_{\pi \in \mathfrak{S}_m} \pi S$$

Interpreting a tableau $S : \lambda \to \{1, \ldots, m\}$ as an element in $\{\lambda\}$ this can be seen as a map $P_m : \{\lambda\} \to \{\lambda\}^{\mathfrak{S}_m}$, where $\{\lambda\}^{\mathfrak{S}_m}$ consists of the elements in $\{\lambda\}$ that are invariant under the action of \mathfrak{S}_m . For example, using the Grassmann-Plücker relations we get that:



Now we consider tableaux that have δ many symbols. Let $\varphi : \{1, \ldots, \delta\} \to \{1, \ldots, m\}$ be a map. For a tableau $T : \lambda \to \{1, \ldots, \delta\}$, we define φT as the tableau that is the result of replacing the entries from $\{1, \ldots, \delta\}$ in T with entries from $\{1, \ldots, m\}$. Let

$$\mathcal{M}_{\delta,m} := \{ \varphi \mid \varphi : \{1, \dots, \delta\} \to \{1, \dots, m\} \text{ is a map} \}.$$

We use this to define

$$M_{\delta,m}T := \sum_{\varphi \in \mathcal{M}_{\delta,m}} \varphi T \in \{\lambda\}^{\mathfrak{S}_m}$$

8 Tableau contraction

Let $\mathbb{S} \in \{\lambda\}$ be the superstandard tableau. Let $\gamma \in \{\lambda\}^*$ denote the vector dual to \mathbb{S} , i.e., the linear map that satisfies $\gamma(\mathbb{S}) = 1$ and $\gamma(T) = 0$ for every semistandard tableau $T \neq \mathbb{S}$.

In the following our goal is to understand γ explicitly in terms of determinants, see eq. (8.2) below. For a matrix g let $g_{1..j,i} \in \mathbb{C}^j$ denote the vector that consists of the top j elements in the *i*th column of g. We interpret a list of j vectors in \mathbb{C}^i as an $i \times j$ matrix of column vectors. A determinant of a matrix with more rows than columns is defined as the determinant of the square matrix of its top rows. For a tableau $T : \lambda \to \{1, \ldots, m\}$ let $T(r, c) \in \{1, \ldots, m\}$ denote the number in row r and column c.

8.1 Lemma. For $g \in \mathbb{C}^{m \times m}$ we have

$$\gamma(gT) = \prod_{c=1}^{\lambda_1} \det(g_{1..\mu_c,T(1,c)},\dots,g_{1..\mu_c,T(\mu_c,c)}),$$
(8.2)

where $\mu = \lambda^t$.

Proof.

$$gT = \sum_{\substack{S:\lambda \to \{1,\dots,m\}}} \left(\prod_{(r,c)\in\lambda} g_{S(r,c),T(r,c)} \right) S = \sum_{\substack{S:\lambda \to \{1,\dots,m\}\\S \text{ regular}}} \left(\prod_{(r,c)\in\lambda} g_{S(r,c),T(r,c)} \right) S$$

The fact that S is superstandard implies that $\gamma(S) = 0$ for all S that do not have a permutation of $\{1, \ldots, \mu_i\}$ in column *i* for all $1 \le i \le \lambda_1$. Therefore

$$\gamma(gT) = \sum_{\substack{S:\lambda \to \{1,\dots,m\}\\ S \text{ is a col-perm. of S}}} \left(\prod_{(r,c)\in\lambda} g_{S(r,c),T(r,c)}\right) \gamma(S) = \sum_{\substack{S:\lambda \to \{1,\dots,m\}\\ S \text{ is a col-perm. of S}}} \left(\prod_{(r,c)\in\lambda} g_{S(r,c),T(r,c)}\right) \prod_{\text{column } c} \operatorname{sgn}(S,c),$$

where sgn(S, c) is the sign of the permutation of column c of S. We conclude

$$\gamma(gT) = \prod_{c=1}^{\lambda_1} \sum_{\pi \in \mathfrak{S}_{\mu_c}} \operatorname{sgn}(\pi) \left(\prod_{r=1}^{\mu_c} g_{\pi(r), T(r, c)} \right),$$

which finishes the proof.

We draw three quick corollaries:

8.3 Corollary. $\gamma(gT) = \prod_{c=1}^{\lambda_1} \gamma(gcol_c)$, where col_c is the c-th column of T.

Proof. Obvious from Lemma 8.1.

8.4 Corollary. For a tableau T that consists of a single regular column of m boxes, we have $\gamma(gT) \in \{-\det(g), \det(g)\}$. In particular, if $g \in SL_m$, we have $\gamma(gT) \in \{-1, 1\}$.

Proof. Since T is regular, the entries of T are precisely all numbers $1, \ldots, m$, not necessarily sorted. Thus according to Lemma 8.1, $\gamma(gT) \in \{-\det(g), \det(g)\}$.

8.5 Corollary. If T is not regular, then $\gamma(gT) = 0$, independent of g.

Proof. $\gamma(gT)$ factorizes according to Cor. 8.3. Since T is not regular, there is a column with a repeated entry. Thus, according to Lemma 8.1, one of the factors of $\gamma(gT)$ equals the determinant of a matrix with a repeated column, and hence is zero.

9 Highest weight functions on the orbit

In this section we prove the following theorem.

9.1 Theorem. The vector space $\operatorname{HWV}_{\lambda^*}(\mathbb{C}[Gp]_d)$ decomposes into a direct sum of vector spaces $\operatorname{HWV}_{\lambda^*}(\mathbb{C}[Gp]_d) = \bigoplus_{\rho \vdash_m d} \mathscr{W}_{\rho}$, and each \mathscr{W}_{ρ} is generated by the functions

$$g \mapsto \gamma(gP_mS),$$

where S runs over all semistandard tableaux S of shape λ and content ρD .

Let \mathscr{W}'_{ϱ} denote the linear space spanned by all semistandard tableaux S of shape λ whose sorted content is ϱD . Then \mathscr{W}_{ϱ} is isomorphic to the \mathfrak{S}_m -invariant subspace of \mathscr{W}'_{ϱ} .

Explicit algebraic Peter-Weyl theorem

 $G \subseteq \mathbb{C}^{m \times m}$ is the nonvanishing set of a polynomial and thus G is a variety. It carries a canonical action of $G \times G$ via $(h', h)g := h'gh^{-1}$ for all $g, h, h' \in G$. This action lifts to the coordinate ring $\mathbb{C}[G]$ via the canonical pullback: $((h', h)f)(g) := f(h'^{-1}gh)$ for all $g, h, h' \in G$, $f \in \mathbb{C}[G]$.

9.2 Theorem (Explicit algebraic Peter-Weyl theorem (see e.g. [Lan17, Thm. 8.6.4.3])). As a $G \times G$ -representation we have

$$\mathbb{C}[G] = \bigoplus_{\lambda} \{\lambda\}^* \otimes \{\lambda\}$$

If we embed $G \hookrightarrow G \times G$ via $g \mapsto (g, id)$, then the λ^* -isotypic component of $\mathbb{C}[G]$ is $\{\lambda\}^* \otimes \{\lambda\}$, which is spanned by functions

 $g \mapsto l(gv),$

where $l \in \{\lambda\}^*$ and $v \in \{\lambda\}$.

H-invariants

Recall the definition of the symmetry subgroup $H \leq G$ from Section 3. Again, consider the action of $G \times G$ on G via $(h', h)g := h'gh^{-1}$ for all $g, h, h' \in G$. The action of $G \times G$ lifts to the coordinate ring $\mathbb{C}[G]$ and we denote by $\mathbb{C}[G]^{\vec{H}}$ the linear subspace of right *H*-invariants. $\mathbb{C}[G]^{\vec{H}}$ carries a left *G*-action. There is a bijection $gp \sim gH$ between points in the orbit Gp and left cosets of *H*. This leads to the following explicit *G*-equivariant algebra isomorphism [TY05, 25.4.6, Prop.]:

Taking right *H*-invariants in Theorem 9.2 yields:

$$\mathbb{C}[G]^{\vec{H}} = \bigoplus_{\lambda} \{\lambda\}^* \otimes \{\lambda\}^H$$

and explicitly: the left λ^* -isotypic component of $\mathbb{C}[G]^{\overrightarrow{H}}$ is spanned by functions

$$g \mapsto l(gv),$$

where $l \in \{\lambda\}^*$ and $v \in \{\lambda\}^H$.

Taking left highest weight functions, we see that $\mathrm{HWV}_{\lambda^*}(\mathbb{C}[G]^{\vec{H}}) \simeq \{\lambda\}^H$ and that $\mathrm{HWV}_{\lambda^*}(\mathbb{C}[G]^{\vec{H}})$ is spanned by functions

$$g \mapsto \gamma(gv),$$
 (9.4)

where $v \in \{\lambda\}^H$. In the next subsection we will use the special structure of H to finish the proof of Theorem 9.1.

The coordinate ring of the orbit of the power sum

Considering eq. (9.4), in order to understand $\text{HWV}_{\lambda^*}(\mathbb{C}[G]^{\vec{H}})$ we need to understand the *H*-invariant space $\{\lambda\}^H$. The group *H* is the semidirect product $\{\lambda\}^H = (\{\lambda\}^{\mathbb{Z}_D^m})^{\mathfrak{S}_m}$, which enables us to analyze $\{\lambda\}^H$ via a two-step process by first analyzing $\{\lambda\}^{\mathbb{Z}_D^m}$ and then taking \mathfrak{S}_m -invariants. We consider the basis of $\{\lambda\}$ given by the semistandard tableaux $S : \lambda \to \{1, \ldots, m\}$. Note that for $g \in \mathbb{Z}_D^m$ we have that *S* and *gS* coincide up to rescaling with a *D*th root of unity. Indeed, if any symbol in the tableau *S* does not occur a multiple of *D* times, then *S* vanishes under the symmetrization

$$S \mapsto \frac{1}{D^m} \sum_{g \in \mathbb{Z}_D^m} \pi S.$$

Moreover, each S in which every number appears a multiple of D many times is fixed under this symmetrization map. Hence $\{\lambda\}^{\mathbb{Z}_D^m}$ has a basis given by semistandard tableaux of shape λ in which each number appears a multiple of D many times. For a partition ϱ let \mathscr{W}'_{ϱ} denote the linear space spanned by all semistandard tableaux of shape λ whose sorted content is ϱD . This gives a decomposition $\{\lambda\}^{\mathbb{Z}_D^m} = \bigoplus_{\varrho} \mathscr{W}'_{\varrho}$ as a direct sum. The action of \mathfrak{S}_m leaves \mathscr{W}'_{ϱ} fixed. Hence $(\{\lambda\}^{\mathbb{Z}_D^m})^{\mathfrak{S}_m} = \bigoplus_{\varrho} (\mathscr{W}'_{\varrho})^{\mathfrak{S}_m}$. A generating set for $\{\lambda\}^H$ is thus obtained by taking the \mathscr{W}'_{ϱ} and their tableau bases and independently symmetrizing over \mathfrak{S}_m , which is done (up to scale) by applying P_m . Using eq. (9.4) we thus obtain the following proposition.

9.5 Proposition. The vector space $\operatorname{HWV}_{\lambda^*}(\mathbb{C}[G]_d^{\vec{H}})$ decomposes into a direct sum of vector spaces $\operatorname{HWV}_{\lambda^*}(\mathbb{C}[G]_d^{\vec{H}}) = \bigoplus_{\varrho \vdash_{md}} \mathscr{W}_{\varrho}$, and each \mathscr{W}_{ϱ} is generated by the functions

$$g \mapsto \gamma(gP_mS),$$

where S runs over all semistandard tableaux S of shape λ and content ρD .

Let \mathscr{W}'_{ϱ} denote the linear space spanned by all semistandard tableaux S of shape λ whose sorted content is ϱD . Then \mathscr{W}_{ϱ} is isomorphic to the \mathfrak{S}_m -invariant subspace of \mathscr{W}'_{ϱ} .

Theorem 9.1 now follows immediately by applying the algebra isomorphism (9.3).

10 A formula for the multiplicities in the coordinate ring of the orbit of the power sum (Proof of Prop. 4.1)

Crucial parts of this section are the result of a collaboration with Greta Panova and appeared in the first author's unpublished lecture notes for a winter 2017/2018 course on geometric complexity theory [Ike19].

In this section we prove Proposition 4.1. In fact, we prove a slightly stronger result as follows.

10.1 Proposition. Let $\{\lambda\}_{\varrho} \subseteq \{\lambda\}$ denote the linear subspace spanned by the tableaux whose sorted content is ϱD . Then dim $(\{\lambda\}_{\varrho})^{\mathfrak{S}_m} = b(\lambda, \varrho, D, d)$.

Proof. Let $\{\lambda\}^{\varrho}$ denote the ϱ -weight space, i.e., the linear space spanned by tableaux of shape λ and content ϱ . Then $\{\lambda\}_{\varrho} = \bigoplus_{\gamma \in \mathfrak{S}_m \rho} \{\lambda\}^{\gamma}$.

For a partition $\rho \vdash_m d$ let $\mathfrak{S}_m \rho \subseteq \mathbb{N}^m$ denote the orbit of ρ . Note that ρ is the only partition in its orbit, while the other lists are not in the correct order. Let stab $\rho \leq \mathfrak{S}_m$ denote the stabilizer of ρ .

10.2 Claim. dim $(\{\lambda\}_{\varrho})^{\mathfrak{S}_m}$ = dim $(\{\lambda\}^{\varrho})^{stab\,\varrho}$.

Proof. We construct an isomorphism of vector spaces.

Let $W := \{\lambda\}$. Let π_1, \ldots, π_r be a system of representatives of left cosets for stab $\varrho \leq \mathfrak{S}_m$ with $\pi_1 = \mathrm{id}$, i.e., $\mathfrak{S}_m = \pi_1 \mathrm{stab} \, \varrho \, \dot{\cup} \, \cdots \, \dot{\cup} \, \pi_r \mathrm{stab} \, \varrho$ and we have $\mathfrak{S}_m \varrho = \{\pi_1 \varrho, \ldots, \pi_r \varrho\}$. Therefore we have the decomposition

$$W_{\varrho} = \bigoplus_{j=1}^{\prime} \pi_j W^{\varrho}.$$

Let $\overline{p}: W_{\varrho} \to W^{\varrho}$ be the projection according to this decomposition. We claim that the restriction

$$p: (W_{\varrho})^{\mathfrak{S}_m} \to (W^{\varrho})^{\mathrm{stab}\,\varrho}$$

is an isomorphism of vector spaces. This then finishes the proof. We verify well-definedness, injectivity, and surjectivity of p.

Well-definedness: The spaces $\pi_1 W^{\varrho}, \ldots, \pi_r W^{\varrho}$ are permuted by \mathfrak{S}_m . Every $\sigma \in \operatorname{stab} \varrho$ fixes W^{ϱ} , thus $\sigma v_1 = v_1$ if $v_1 \in W^{\varrho}$. Thus the map $v = \sum_{j=1}^r v_j \stackrel{\overline{\rho}}{\mapsto} v_1$ maps W_{ϱ} to $(W^{\varrho})^{\operatorname{stab} \varrho}$. Injectivity: If $v \in (W_{\varrho})^{\mathfrak{S}_m}$, then $v = \pi v = \sum_j \pi v_j$. Therefore $v_j = \pi_j v_1$. If p(v) = 0, then

 $v_1 = 0$, thus all $v_j = 0$, which proves injectivity.

Surjectivity: Let $v_1 \in (W^{\varrho})^{\operatorname{stab} \varrho}$. Set $v_j := \pi_j v_1$ and put $v := \sum_i v_j$. Clearly $p(v) = v_1$. It remains to verify that v is \mathfrak{S}_m -invariant.

$$v = \sum_{j=1}^{r} \pi_j v_1 = \sum_{j=1}^{r} \frac{1}{|\operatorname{stab} \varrho|} \sum_{\tau \in \operatorname{stab} \varrho} \pi_j \tau v_1 = \frac{1}{|\operatorname{stab} \varrho|} \sum_{\pi \in \mathfrak{S}_m} \pi v_1,$$

et.

which is \mathfrak{S}_m -invariant.

We are left with determining dim $(\{\lambda\}^{\varrho})^{\operatorname{stab} \varrho}$. We use a detour via Specht modules: the Specht module $[\lambda]$ is an irreducible $\mathfrak{S}_{|\lambda|}$ -representation that can be constructed as the subrepresentation of $\{\lambda\}$ spanned by all standard tableaux. Define the Young subgroup $G_{\varrho} \subseteq \mathfrak{S}_{dD} := \mathfrak{S}_{\varrho_1 D} \times \cdots \times \mathfrak{S}_{\varrho_m D}$. We use that $\{\lambda\}^{\varrho} \simeq [\lambda]^{G_{\varrho}}$, see e.g. [Ike12, Sec. 4.3(A)].

Schur-Weyl duality implies that

$$\dim \left([\lambda]^{G_{\varrho}} \right)^{\operatorname{stab} \varrho} = \dim \operatorname{HWV}_{\lambda}(\{\lambda\} \otimes ([\lambda]^{G_{\varrho}})^{\operatorname{stab} \varrho}) = \dim \operatorname{HWV}_{\lambda}((\otimes^{dD} V)^{G_{\varrho} \rtimes \operatorname{stab} \varrho})$$

for V having large enough dimension.

$$(\otimes^{dD}V)^{G_{\varrho} \rtimes \operatorname{stab} \varrho} = (\operatorname{Poly}^{D\varrho_{1}}V \otimes \cdots \otimes \operatorname{Poly}^{D\varrho_{m}}V)^{\operatorname{stab} \varrho}$$

$$= \left(\bigotimes^{\hat{\varrho}_{1}} \operatorname{Poly}^{D}V \otimes \bigotimes^{\hat{\varrho}_{2}} \operatorname{Poly}^{2D}V \otimes \cdots \otimes \bigotimes^{\hat{\varrho}_{d}} \operatorname{Poly}^{dD}V \right)^{\operatorname{stab} \varrho}$$

$$= \underbrace{\mathbb{C}[\operatorname{Poly}^{D}V]_{\hat{\varrho}_{1}}^{*}}_{=\bigoplus_{\mu^{1}}\{\mu^{1}\}^{\oplus a_{\mu^{1}}(\hat{\varrho}_{1},D)}} \otimes \mathbb{C}[\operatorname{Poly}^{2D}V]_{\hat{\varrho}_{2}}^{*} \otimes \cdots \otimes \underbrace{\mathbb{C}[\operatorname{Poly}^{dD}V]_{\hat{\varrho}_{d}}^{*}}_{=\bigoplus_{\mu^{d}}\{\mu^{d}\}^{\oplus a_{\mu^{d}}(\hat{\varrho}_{d},dD)}} (\dagger).$$

The multiplicity of $\{\mu^i\}^*$ in $\mathbb{C}[\operatorname{Poly}^{iD}V]_{\hat{\varrho}_i}$ is $a_{\mu^i}(\hat{\varrho}_i, iD)$. Using distributivity we obtain that the multiplicity of $\{\lambda\}$ in the representation (†) equals

$$\sum_{\substack{\mu^{1},\mu^{2},...,\mu^{d}\\\mu^{i}\vdash\hat{\varrho}_{i}Di}} c_{\mu^{1},\mu^{2},...,\mu^{d}} \prod_{i=1}^{u} a_{\mu^{i}}(\hat{\varrho}_{i},iD)$$

Proof of Proposition 4.1.
$$\operatorname{mult}_{\lambda} \mathbb{C}[Gp]_d \stackrel{\operatorname{Thm. }9.1}{=} \sum_{\varrho \vdash_m d} (\mathscr{W}_{\varrho})^{\mathfrak{S}_m} \stackrel{\operatorname{Prop. }10.1}{=} \sum_{\varrho \vdash_m d} b(\lambda, \varrho, D, d).$$

11 Highest weight functions on the orbit closure

In this section we study the coordinate ring $\mathbb{C}[\overline{Gp}]$. Much less is known about this ring compared to $\mathbb{C}[Gp]$, in particular we do not have formulas for $\mathbb{C}[\overline{Gp}]$ that are comparable to Prop. 4.1. Nevertheless, the following theorem provides a way to analyze $\mathbb{C}[\overline{Gp}]$ and connect it to $\mathbb{C}[Gp]$.

For an $m \times n$ matrix A let $Ap := \ell_1^D + \ell_2^D + \cdots + \ell_n^D$, where ℓ_i is the linear form given by the *i*-th column of A. Note that this is a generalization of gp for $g \in GL_m$.

11.1 Theorem. If $|\lambda|$ is not divisible by D, then $\text{HWV}_{\lambda^*}(\mathbb{C}[\overline{Gp}]) = 0$. Let $\lambda \vdash \delta D$. The vector space $\text{HWV}_{\lambda^*}(\mathbb{C}[\overline{Gp}])$ is generated by the functions

$$g \mapsto \gamma(gM_{\delta,m}T),$$

where T runs over all semistandard tableaux of shape λ in which each entry $1, \ldots, \delta$ appears exactly D many times.

Additionally, for a semistandard tableau T of shape λ in which each entry $1, \ldots, \delta$ appears exactly D many times, the function

$$A \mapsto \gamma(AM_{\delta,m}T)$$

is either zero or a HWV of weight λ^* in $\mathbb{C}[\operatorname{Poly}^D \mathbb{C}^m]$.

This is a rephrasing of [AIR16], which is a special case of [BIP19, Prop. 4.5 and Thm. 4.7]. Indeed, [BIP19] covers more general cases. For the sake of completeness and to highlight that the proof technique is very different from the technique in Section 9, we prove Theorem 11.1.

Proof. The first observation follows from the fact that $\overline{Gp} \subseteq \text{Poly}^D \mathbb{C}^m$ is a closed subvariety that is closed under rescaling, and hence $\mathbb{C}[\overline{Gp}]$ is a graded subalgebra of $\mathbb{C}[\text{Poly}^D \mathbb{C}^m]$. We know that in each degree δ component $\mathbb{C}[\text{Poly}^D \mathbb{C}^m]_{\delta}$ of $\mathbb{C}[\text{Poly}^D \mathbb{C}^m]$ the only types λ^* that occur satisfy $\lambda \vdash \delta D$, see e.g. [Ike12, Lemma 4.3.3].

Schur-Weyl duality yields that

$$\bigotimes^{\delta D} \mathbb{C}^{m*} = \bigoplus_{\lambda \vdash_m \delta D} \{\lambda^*\} \otimes [\lambda].$$
(11.2)

A highest weight vector of weight λ^* in $\bigotimes^{\delta D} \mathbb{C}^{m*}$ is given for example by

$$v_{\lambda} := x_1 \wedge x_2 \wedge \cdots \wedge x_{\mu_1} \otimes x_1 \wedge x_2 \wedge \cdots \wedge x_{\mu_2} \otimes \cdots \otimes x_1 \wedge x_2 \wedge \cdots \wedge x_{\mu_{\lambda_1}},$$

where $\mu = \lambda^t$ and $\{x_i\}_i$ is the basis of \mathbb{C}^{m*} . Let T_{λ} denote the column-standard tableau of shape λ , i.e., the tableau that is filled with the numbers $1, \ldots, |\lambda|$ in a columnwise fashion from left to right, top to bottom. Since $[\lambda]$ is irreducible, from (11.2) we see that $\text{HWV}_{\lambda^*}(\bigotimes^{\delta D} \mathbb{C}^{m*})$ is generated by the set $\{v_{\lambda}\pi \mid \pi \in \mathfrak{S}_{\delta D} \text{ s.t. } \pi T_{\lambda} \text{ is standard}\}.$

By the polarization principle (see e.g. [Ike12, Claim 4.2.13]), all functions f in $\mathbb{C}[\operatorname{Poly}^D \mathbb{C}^{m*}]_{\delta}$ can be obtained via some tensor $v_f \in \bigotimes^{\delta D} \mathbb{C}^m$ and defining

$$f(y) := \langle v_f, y^{\otimes \delta} \rangle.$$

The resulting function f is a highest weight function iff v_f is a HWV. Thus we see that $\text{HWV}_{\lambda^*}(\mathbb{C}[\text{Poly}^D\mathbb{C}^m]_{\delta})$ is generated by the functions

$$f(y) := \langle \pi v_{\lambda}, y^{\otimes \delta} \rangle.$$

In the following we analyze how to restrict these functions to \overline{Gp} . When $y = gp = \ell_1^D + \cdots + \ell_m^D$ is in the orbit of the power sum, then clearly

$$y^{\otimes \delta} = \sum_{\varphi:\{1,\dots,\delta\} \to \{1,\dots,m\}} \ell^D_{\varphi(1)} \otimes \dots \otimes \ell^D_{\varphi(\delta)}.$$

The linear forms ℓ_i correspond to the vectors $g_{1..m,i}$. We evaluate

$$f(y) := \langle \pi v_{\lambda}, y^{\otimes \delta} \rangle = \langle v_{\lambda}, \pi(y^{\otimes \delta}) \rangle$$
$$= \sum_{\substack{\beta:\{1,\dots,\delta D\} \to \{1,\dots,m\} \\ \text{respecting } \pi T_{\lambda}}} \prod_{c=1}^{\lambda_{1}} \det(g_{1..\mu_{c},\beta(\pi T_{\lambda}(1,c))},\dots,g_{1..\mu_{c},\beta(\pi T_{\lambda}(\mu_{c},c))}),$$

where β respects a tableau S if all numbers $1, \ldots, D$ are mapped to the same value, and all numbers $D + 1, \ldots, 2D$ are mapped to the same value, and so on.

Therefore the vector space $HWV_{\lambda^*}(\mathbb{C}[\overline{Gp}])$ is generated by the functions

$$g \mapsto \sum_{\substack{\beta:\{1,\dots,Dd\} \to \{1,\dots,m\}\\ \text{respecting } S}} \prod_{c=1}^{\lambda_1} \det(g_{1..\mu_c,\beta(S(1,c))},\dots,g_{1..\mu_c,\beta(S(\mu_c,c))}),$$
(11.3)

where S runs over all standard tableaux of shape λ .

Given a standard tableau S we define a tableau T by replacing the first D entries $1, \ldots, D$ by the number 1, the next D entries $D+1, \ldots, 2D$ by the number 2, and so on. It is easy to check that if T is not regular, then the function corresponding to S describes the zero function, because each summand in (11.3) has a zero factor that is the determinant of a matrix with a repeating column. We assume from now on that T is regular. Since T is regular and S is standard, T is semistandard. We rewrite (11.3) as follows:

$$g \mapsto \sum_{\varphi:\{1,\dots,\delta\}\to\{1,\dots,m\}} \prod_{c=1}^{\lambda_1} \det(g_{1..\mu_c,\varphi(T(1,c))},\dots,g_{1..\mu_c,\varphi(T(\mu_c,c))})$$
(11.4)

Using (8.2), we can write this in terms of γ as follows:

$$g\mapsto \sum_{\varphi:\{1,\ldots,\delta\}\to\{1,\ldots,m\}}\gamma(g\varphi T).$$

By definition of $M_{\delta,m}$, this can be rewritten as:

$$g \mapsto \gamma(gM_{\delta,m}T)$$

which finishes the proof of the second part of Theorem 11.1.

For the last claim, we note that in this construction of highest weight functions we did not use that g is a square matrix. A rectangular matrix A works in the same way.

12 An equation for Waring rank (Proof of the missing part in Prop. 5.3)

In this section we prove the following proposition that was used in the proof of Proposition 5.3.

12.1 Proposition. Let $\nu = (2m) + m \times 2m$. Then $a_{\nu}(2(m+1), m) > 2$.

Proof. By Proposition 4.1 and Lemma 5.2 we have $\operatorname{mult}_{\nu^*}\mathbb{C}[\overline{Gp}] = \operatorname{mult}_{\nu^*}\mathbb{C}[Gp] = 2$. This means that there are two linearly independent HWVs of weight ν^* that do not vanish on Gp. To finish the proof it suffices to construct a nonzero HWV of weight ν^* that vanishes on Gp, because then these three HWVs are linearly independent. Note that in particular we construct an equation that vanishes on all polynomials of Waring rank at most m.

We use the last part of Theorem 11.1 to construct this third function. Let n := 2m + 2. Let T_{left} be the $m \times (m + 2)$ tableau that is filled in a rowwise fashion from top to bottom and from left to right with m many 1s, then m many 2s, and so on, until m many (m + 2)s. For example, if m = 6, then

$T_{left} =$	1	1	1	1	1	1	2	2
- Iert	2	2	2	2	3	3	3	3
	3	3	4	4	4	4	4	4
	5	5	5	5	5	5	6	6
	6	6	6	6	7	7	7	7
	7	7	8	8	8	8	8	8

We remark that the function corresponding to T_{left} via Theorem 11.1 is a generalization of Aronhold's degree 4 invariant on ternary cubics, see also [BI17] for other (related) generalizations.

Let T_{right} be the $(m \times 2) + (m^2 - 2m)$ tableau whose first two columns are equal and consist of the entries $m + 3, m + 4, \ldots, 2m + 2$ from top to bottom. The remaining singleton columns get filled with m - 2 many entries m + 3, m - 2 many entries m + 4, m - 2 many entries m + 5, and so on, until m - 2 many entries 2m + 2. For example, if m = 6, then

$T_{right} =$	9	9	9	9	9	9	10	10	10	10	11	11	11	11	12	12	12	12	13	13	13	13	14	141	414
ingine	10	10																							
	11	11																							
	12	12																							
	13	13																							
	14	14																							

The tableau T is defined as the concatenation $T := T_{\text{left}} + T_{\text{right}}$. We observe that T is duplex.

By Theorem 11.1 the function $f : A \mapsto \gamma(AM_{n,n}T)$ is either zero or a HWV of weight ν^* in $\mathbb{C}[\operatorname{Poly}^m \mathbb{C}^m]$.

12.2 Claim. f does not vanish on $Poly^m \mathbb{C}^m$.

Proof. This is due to the fact that T is duplex, in complete analogy to [BCI11]. Choose A to be an $m \times n$ matrix whose entries are real numbers chosen generically (one can alternatively think of the entries being chosen uniformly at random for example from a Gaussian distribution). Since Tis duplex, each summand in $\gamma(AM_{n,n}T)$ is a product of determinants (see (11.4)), but each factor in the product appears an even number of times and hence the product is nonnegative. Since Awas chosen generically, for the identity map id $\in \mathcal{M}_{n,n}$ we have $\gamma(A \operatorname{id} T) > 0$. Any finite sum of nonnegative numbers that contains at least one positive number is nonzero, so $\gamma(AM_{n,n}T)$ is nonzero. This finishes the proof.

The preceding claim implies that f is a nonzero HWV of weight ν^* in $\mathbb{C}[\operatorname{Poly}^m \mathbb{C}^m]$. To finish the proof of Proposition 12.1 it suffices to prove that f vanishes on Gp. The crucial property is that no tableau in $\mathcal{M}_{m+2,m}T_{\mathsf{left}}$ is regular: Since T_{left} is rectangular with the maximum number of rows, a regular tableau in $\mathcal{M}_{m+2,m}T_{\mathsf{left}}$ has m+2 many 1s, m+2 many 2s, and so on, but every symbol in $\mathcal{M}_{m+2,m}T_{\mathsf{left}}$ appears a multiple of m many times. Since no tableau in $\mathcal{M}_{m+2,m}T_{\mathsf{left}}$ is regular, no tableau in $\mathcal{M}_{n,m}T$ is regular. Hence all summands in $\gamma(gM_{n,m}T)$ are zero, see Corollary 8.5. This finishes the proof of Proposition 12.1.

13 Proof of the Main Technical Theorem 4.2 using the Tableau Lifting Theorem

In this section we prove Theorem 4.2, based on the following combinatorial Tableau Lifting Theorem 13.1 whose long and combinatorial proof we will develop during the remaining sections of this paper. Much simpler forms of other tableau lifting theorems appeared in [KL14, BIP19].

Fix a shape λ and natural numbers m and e. For a tableau T of shape $(m \times e) + \lambda$ we define leftpart(T) to be the $m \times e$ rectangular subtableau consisting of the leftmost e columns, and we define rightpart(T) to be the shape λ subtableau consisting of the rightmost λ_1 columns. In particular, we have T = leftpart(T) + rightpart(T). For $\varrho \in (\mathbb{N}_{\geq 0})^m$ a tableau S has content $D\varrho$ if each number iappears exactly $D\varrho_i$ many times in S.

13.1 Theorem (Tableau Lifting Theorem). Let $D \ge 3$ and $m \ge 2$. Given a regular tableau S of shape $\lambda \vdash_m dD$ and content $D\varrho$ for $\varrho \vdash_m d$. Let

- $e_{\varrho} := \sum_{i=1}^{m} \lceil \frac{\varrho_i}{D-2} \rceil$ in the case where D is even
- and $e_{\varrho} := \sum_{i=1}^{m} 2 \lceil \frac{\varrho_i}{2(D-2)} \rceil$ in the case where D is odd and $\binom{2(D-1)}{D-1} \ge 2(m-1)$.

Let $\delta := me_{\varrho} + d$. In both cases there exists a tableau $T : \lambda + (m \times e_{\varrho}D) \to \{1, \ldots, \delta\}$ in which every entry appears exactly D many times such that

- 1. For each $\varphi \in \mathcal{M}_{\delta,m}$ for which $\varphi(T)$ is regular we have that $\operatorname{rightpart}(\varphi(T)) \in \mathfrak{S}_m S$,
- 2. For each $\varphi \in \mathcal{M}_{\delta,m}$ for which $\varphi(T)$ is regular we have that $\mathsf{leftpart}(\varphi(T))$ is duplex,
- 3. there exists $\varphi \in \mathcal{M}_{\delta,m}$ such that $\varphi(T)$ is regular and rightpart $(\varphi(T)) = S$.

Proof of the Main Theorem 4.2. By Theorem 9.1 and Proposition 10.1 we know that there exists a set of regular tableaux $\{S_{\varrho,i} \mid \varrho \in \Xi, 1 \leq i \leq b(\lambda, \varrho, D, d)\}$ of shape λ such that each $S_{\varrho,i}$ has content $D\varrho$ and the set of corresponding functions

$$f^{S_{\varrho,i}} \in \mathbb{C}[Gp], \quad gp \mapsto \gamma(gP_m S_{\varrho,i}).$$

is linearly independent. Since all these functions $f^{S_{\varrho,i}}$ are homogeneous of the same degree d, they are not only linearly independent as functions on $\operatorname{GL}_m p$, but also their restrictions to $\operatorname{SL}_m p$ are linearly independent. Using the Tableau Lifting Theorem 13.1, for each $S_{\varrho,i}$ we construct a tableau $T_{\varrho,i}$ of shape $\lambda + (m \times e_{\varrho} D)$ satisfying the properties listed in Theorem 13.1. We claim that for all ϱ, i there exists $\alpha \neq 0$ such that

under the map $\psi : \varphi \mapsto \operatorname{rightpart}(\varphi T_{\varrho,i})$ that maps from $\mathcal{M}_{\delta,m}$ to the set of tableaux of shape λ , each tableau in $\mathfrak{S}_m S_{\varrho,i}$ has exactly α many preimages in $\mathcal{M}_{\delta,m}$ for which (*) $\varphi T_{\varrho,i}$ is regular.

Proof of (*): Clearly $\varphi T_{\varrho,i}$ is regular iff $\pi \varphi T_{\varrho,i}$ is regular. Note that ψ is \mathfrak{S}_m -equivariant in the following sense: $\psi(\pi \circ \varphi) = \pi \psi(\varphi)$. Hence taking the preimage ψ^{-1} is also \mathfrak{S}_m -equivariant. Thus for all $\hat{S} \in \mathfrak{S}_m S_{\varrho,i}$ we have $\varphi \in \psi^{-1}(\hat{S})$ iff $\pi \circ \varphi \in \psi^{-1}(\pi \hat{S})$. Thus the application of π gives a

bijection between the preimages of \hat{S} and $\pi \hat{S}$. To prove claim (*) it remains to show that $\alpha \neq 0$. This follows from Thm. 13.1(3). \Box

According to [BI17, Prop. 3.25] there exists an SL_m -invariant function Φ in $\mathbb{C}[\overline{Gp}]$ with $\Phi(p) = 1$, called the fundamental invariant. Moreover,

$$\begin{cases} \Phi \text{ has degree } m \text{ and } \Phi(gp) = \det(g)^D p \text{ for } g \in \mathrm{GL}_m & \text{if } D \text{ is even} \\ \Phi \text{ has degree } 2m \text{ and } \Phi(gp) = \det(g)^{2D} p \text{ for } g \in \mathrm{GL}_m & \text{if } D \text{ is odd and } 2m \leq {2D \choose D}. \end{cases}$$

Since orbit closures are irreducible varieties, given two highest weight vectors f of weight λ and \tilde{f} of weight $\tilde{\lambda}$ in $\mathbb{C}[\overline{Gp}]$, their product $f \cdot \tilde{f}$ is nonzero and a highest weight vector of weight $\lambda + \tilde{\lambda}$. Let $\overline{f}^{S_{\varrho,i}}$ be the product of $\Phi^{e_{\Xi}-e_{\varrho}}$ ($\Phi^{(e_{\Xi}-e_{\varrho})/2}$ if D is odd) and the function corresponding to $T_{\varrho,i}$:

$$\overline{f}^{S_{\varrho,i}}(gp) = \det(g)^{D(e_{\Xi}-e_{\varrho})} \cdot \gamma(gM_{m,d}T_{\varrho,i}) \qquad \text{(holds for } D \text{ even and odd.)}$$

We claim that $\overline{f}^{S_{\varrho,i}}$ coincides with $f_{\varrho,i}^S$ when restricted to $\mathrm{SL}_m p$ (up to nonzero a factor), which can be seen as follows. For $g \in \mathrm{SL}_m$ we have

$$\begin{split} \overline{f}^{S_{\varrho,i}}(gp) &= 1 \cdot \gamma(gM_{m,d}T_{\varrho,i}) = \sum_{\varphi \in \mathcal{M}_{m,d}} \gamma(g\varphi T_{\varrho,i}) \\ &\stackrel{\text{Cor. 8.5}}{=} \sum_{\substack{\varphi \in \mathcal{M}_{m,d} \\ \varphi T_{\varrho,i} \text{ regular}}} \gamma(g\varphi T_{\varrho,i}) \\ &\stackrel{\text{Cor. 8.3 and Cor. 8.4 and Thm. 13.1(2)}}{=} (\pm 1)^2 \sum_{\substack{\varphi \in \mathcal{M}_{m,d} \\ \varphi T_{\varrho,i} \text{ regular}}} \gamma(grightpart(\varphi T_{\varrho,i})) \\ &\stackrel{\text{Thm. 13.1(1) and (*)}}{=} \alpha \sum_{\substack{\hat{S} \in \mathfrak{S}_m S_{\varrho,i} \\ g \in \mathcal{M}_{p,i} \in \mathfrak{S}_{p,i} \in \mathfrak{S}_{p,i}}} \gamma(g\hat{S}) = \frac{|\mathfrak{S}_m S_{\varrho,i}|\alpha}{m!} \gamma(gP_m S_{\varrho,i}) \\ &= \frac{|\mathfrak{S}_m S_{\varrho,i}|\alpha}{m!} f^{S_{\varrho,i}}(gp). \end{split}$$

Since each $\overline{f}^{S_{\varrho,i}}$ coincides with $f^{S_{\varrho,i}}$ when restricted to $\operatorname{SL}_m p$ (up to a nonzero factor), the $\overline{f}^{S_{\varrho,i}}$ are linearly independent. It follows from Theorem 11.1 that the $\overline{f}^{S_{\varrho,i}}$ are restrictions of functions in $\operatorname{HWV}_{(\lambda+(m\times e_{\Xi}D))^*}\mathbb{C}[\overline{Gp}]$. This proves the main Theorem 4.2.

It remains to prove the Tableau Lifting Theorem 13.1, whose purely combinatorial proof will be the focus of the rest of the paper. The construction for odd D is more complicated than for even D, which is why we focus on the case where D is even first.

We will construct $\mathsf{leftpart}(T)$ and $\mathsf{rightpart}(T)$ mainly independently.

For even D, the alphabet that we are using for T is not $\{1, \ldots, \delta\}$, but a more descriptive alphabet using the symbols i_{ℓ} and j_k^i . For each box \Box in S, if \Box has the entry i, then the box corresponding to \Box in rightpart(T) has the symbol i_{ℓ} for some ℓ . The only other constraints for rightpart(T) are concerned with how often the different symbols i_{ℓ} appear. The symbols j_k^i do not appear in rightpart(T), but only in leftpart(T). The tableau leftpart(T) is constructed in several steps, starting with a tableau obtained from a set of hypergraphs $H^{(i)}$, and then reordering entries within the rows.

For odd D the situation is similar. The alphabet that we are using for T is not $\{1, \ldots, \delta\}$, but a more descriptive alphabet using the symbols i_{ℓ}, j_k^i, j_k^i . For each box \Box in S, if \Box has the entry i, then

the box corresponding to \Box in rightpart(T) has the symbol i_{ℓ} for some ℓ . The only other constraints for rightpart(T) are concerned with how often the different symbols i_{ℓ} appear. Symbols j_k^i and j_k^i do not appear in rightpart(T), but only in leftpart(T). The tableau leftpart(T) is constructed in several steps, starting with a tableau obtained from a set of hypergraphs $H^{(i)}$ (similar to those hypergraphs in the case where D is even), and then reordering entries within the rows.

The construction for odd D has many more subtleties than the construction for even D and there are numerous slight differences between the two cases. We think that the readability would suffer greatly if we did not explain the whole construction again for odd D in a self-contained manner, including the parts that are very similar to the even case. Therefore in the following sections we first treat the case for even D and then treat the case for odd D in a fairly self-contained manner. The reader will see that much more care and attention to the details is necessary in the case where D is odd.

14 The hypergraphs $H^{(i)}$ for even D

Let $I := \{i \mid \varrho_i \neq 0\}$. In order to construct $\mathsf{leftpart}(T)$ we will first construct a so-called (D, ϱ_i) -hypergraph for each $i \in I$. We refer to that hypergraph as $H^{(i)}$. The number of columns in $\mathsf{leftpart}(T)$ will precisely be the number of vertices in all these hypergraphs together. So we want the hypergraphs to be as small as possible.

We first recall some basic terms. Let H = (V, E) by a hypergraph and $e \in E$ be an edge. Then we define the *size* of an edge size(e) as the number of vertices in e. Let $v, w \in V$. Then a *path* between v and w is a sequence of edges (e_1, e_2, \ldots, e_l) such that $v \in e_1$, $w \in e_l$, and $e_i \cap e_{i+1} \neq \emptyset$. We say that two vertices are connected in H iff there exists a path between them. We say that a hypergraph is *connected* iff every pair of vertices is connected. For a nonempty set X a set partition P of X is a set of pairwise disjoint subsets whose union is X.

14.1 Definition. Let D, K be integers. A (D, K)-hypergraph is defined to be a hypergraph H = (V, E) that satisfies the following properties:

- (1) H is connected.
- (2) H has two different types of hyperedges: the block edges and the name edges.
- (3) Each block edge has size D, and the set of block edges $E_{\text{Block}} \subseteq E$ is a set partition of V.
- (4) Each name edge has size strictly less than D, but at least size 1, and the set of name edges $E_{\text{Name}} \subseteq E$ is a set partition of V.
- (5) $|E_{\text{Name}}| |E_{\text{Block}}| = K.$
- (6) There exists a name edge e_{Name} and a block edge e_{Block} whose intersection contains at least 2 vertices. We choose one of these two vertices and call it the *link vertex*.

Several examples for (6, K)-hypergraphs are given in Figure 1.

14.2 Proposition. For even $D \ge 4$, $K \ne 0$, there exists a (D, K)-hypergraph that has exactly $\lceil \frac{K}{D-2} \rceil$ many block edges.

Proof. Let $n := \lceil \frac{K}{(D-2)} \rceil$. We start the construction by considering n many disjoint block edges with D many vertices each. We arrange the vertices in a linear fashion as in Figure 1. The leftmost vertex is the link vertex. We now place the vertices in K + n many name edges as follows. As in

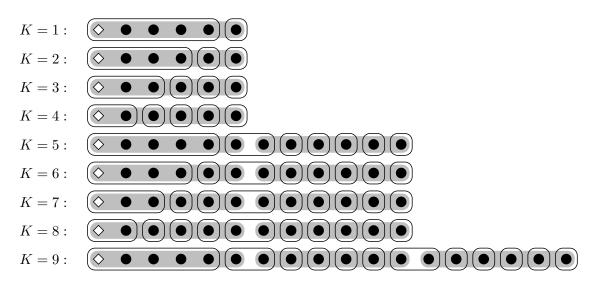


Figure 1: (6, K)-hypergraphs for several values of K. Each block edges is represented as a filled gray blob. Each name edges is enclosed by a solid black curve. The link vertices are drawn as white diamonds.

Figure 1, the rightmost vertex of every block edge but the last shall be placed in a size 2 name edge with the leftmost vertex of the next block edge. The resulting hypergraph is connected. At this point we have nD - 2(n-1) = n(D-2) + 2 vertices that are not in name edges yet; and we have K + n - (n-1) = K + 1 name edges left to put vertices in. Since

$$(n(D-2)+2) - (K+1) = \left(\left\lceil \frac{K}{(D-2)} \right\rceil (D-2) + 2 \right) - (K+1) \ge (K+2) - (K+1) = 1 > 0,$$

we can position the name edges so that the link vertex is in a name edge of size at least 2. Moreover, we position that name edge of size at least 2 in such a way that the link vertex has a vertex that not only lies in the same name edge, but also in the same block edge. \Box

15 Construction of $\mathsf{leftpart}(T)$ for even D

For each $i \in I$ let $H^{(i)}$ be a (D, ϱ_i) -hypergraph from Proposition 14.2. We write $E_{\text{Block}}^{(i)}$ to denote its set of block edges and $E_{\text{Name}}^{(i)}$ to denote its set of name edges.

In this section, for every $i \in I$ and every $e \in E_{\text{Block}}^{(i)}$ we construct an $m \times D$ block tableau \check{B}_e such that $\mathsf{leftpart}(T)$ is constructed as the concatenation

$$\mathsf{leftpart}(T) := \sum_{i \in I} \sum_{e \in E_{\mathrm{Block}}^{(i)}} \check{B}_e.$$
(15.1)

Notice that since every block edge has size D (see Def. 14.1((3))), this implies that the number of columns in leftpart(T) is equal to the total number of vertices in the hypergraphs $H^{(i)}$, $i \in I$.

Each $m \times D$ block tableau \check{B}_e is constructed in three steps: First we construct an $m \times D$ block tableau B_e in which each column corresponds to a vertex in $H^{(i)}$, then we exchange entries between columns that correspond to link vertices.

Let $\zeta^{(i)}$ denote the link vertex in $H^{(i)}$. We attach some additional data to each $H^{(i)}$ as follows. We put a linear order on the set of name edges $E_{\text{name}}^{(i)}$ and for each vertex v in $H^{(i)}$ we define $\ell(v)$ to

\diamond				\bigcirc			\bigcirc	\bigcirc	\bigcirc		
1_{1}^{2}											
2_1	2_1	2_1	2_1	2_2	2_3	2_3	2_4	2_5	2_6	2_7	2_8
3_{1}^{2}	3_1^2	3_{1}^{2}	3_1^2	3_{1}^{2}	3_{1}^{2}	3_{2}^{2}	3_{2}^{2}	3_{2}^{2}	3_{2}^{2}	3_{2}^{2}	3_{2}^{2}
4_1^2	4^2_1	4^{2}_{1}	4_1^2	4_1^2	4^{2}_{1}	4^2_2	4^2_2	4^2_2	4^2_2	4^2_2	4^2_2

Figure 2: Here i = 2. A (6,6)-hypergraph $H^{(2)}$ with two block edges e_1 and e_2 and the corresponding concatenated tableau $B_{e_1} + B_{e_2}$. Vertices are drawn directly above their corresponding columns. To make the value of i easy to see, the second row is highlighted.

be the index of its corresponding name edge. Here $\ell(v) = 1$ if v lies in the first name edge, $\ell(v) = 2$ for the next name edge, and so on. We ensure that

$$\ell(\zeta^{(i)}) = 1. \tag{15.2}$$

In the same way, we put a linear order on the set of block edges; for each block edge e we write k(e) for its index and for each vertex v in $H^{(i)}$ we define k(v) to be the index of its corresponding block edge. We ensure that

$$k(\zeta^{(i)}) = 1. \tag{15.3}$$

Moreover, for every vertex v in any $H^{(i)}$ we define i(v) := i.

In the following, for each vertex v we define an $m \times 1$ rectangular tableau (i.e., a column of length m) called B_v . Concatenating them results in $B_e := \sum_{v \in e} B_v$. The order of columns does not matter, but it is convenient to have the vertices of $H^{(i)}$ ordered from left to right in the same way as the columns of $\mathsf{leftpart}(T)$. Later we define \check{B}_e , from which we can extract \check{B}_v for $v \in e$ as follows: If B_v is the *n*-th column of B_e , then \check{B}_v is the *n*-th column of \check{B}_e .

Starting with B

Let e be in the k-th block edge in $H^{(i)}$ and let $v \in e$. The column B_v is defined by the following properties.

the *i*-th entry of
$$B_v$$
 is $i_{\ell(v)}$ (15.4)

the *j*-th entry
$$(j \neq i)$$
 of B_v is j_k^i (15.5)

An example is given in Figure 2.

From B to \check{B}

Most columns B_v and \check{B}_v coincide, as we define $\check{B}_v := B_v$ if $v \notin \{\zeta^{(i)} \mid i \in I\}$. This means that only the columns corresponding to link vertices are adjusted.

Let *h* denote the smallest number in *I*. For $i \in I$, $i \neq h$, the column $\dot{B}_{\zeta^{(i)}}$ arises from $B_{\zeta^{(i)}}$ by switching the *i*-th entry with the *i*-th entry in $B_{\zeta^{(h)}}$. This means that the column $\check{B}_{\zeta^{(h)}}$ arises from $B_{\zeta^{(i)}}$ by switching the *i*-th entry with the *i*-th entry in $B_{\zeta^{(i)}}$ for all $i \in I$.

The columnwise description of B_v thus as follows.

\diamond						\bullet	\bigcirc	\bigcirc	\bigcirc	\bigcirc	\bullet	\diamond						\diamond		•	•	•)(
1_1	1_{1}	1_{1}	1_{1}	1_{1}	1_{2}	1_{2}	1_3	1_4	1_5	1_6	1_7	1_{1}^{2}	1_{1}^{2}	1_{1}^{2}	1_{1}^{2}	1_{1}^{2}	1_{1}^{2}	1^{3}_{1}	1^{3}_{1}	1^{3}_{1}	1_{1}^{3}	1_{1}^{3}	1^{3}_{1}
2^{1}_{1}	2^{1}_{1}	2^{1}_{1}	2^{1}_{1}	2^{1}_{1}	2^{1}_{1}	2^{1}_{2}	2^{1}_{2}	2^{1}_{2}	2^{1}_{2}	2^{1}_{2}	2^{1}_{2}	2_1	2_1	2_1	2_1	2_{2}	2_3	2^{3}_{1}	2^{3}_{1}	2^{3}_{1}	2^{3}_{1}	2^{3}_{1}	2^{3}_{1}
3^{1}_{1}	3^{1}_{1}	3_{1}^{1}	3^{1}_{1}	3_{1}^{1}	3^{1}_{1}	3^{1}_{2}	3^{1}_{2}	3^{1}_{2}	3^{1}_{2}	3^{1}_{2}	3^{1}_{2}	3_{1}^{2}				3_1^2	-	-	3_1	3_1	3_1	3_1	3_2
4_{1}^{1}	4^{1}_{1}	4^{1}_{1}	4^{1}_{1}	4^{1}_{1}	4^{1}_{1}	4^{1}_{2}	4^{1}_{2}	4^{1}_{2}	4^{1}_{2}	4^{1}_{2}	4^{1}_{2}	4_1^2	4^{2}_{1}	4_1^2	4_1^2	4_1^2	4_1^2	4^{3}_{1}	4_1^3	4^{3}_{1}	4_1^3	4^{3}_{1}	4_1^3
1_1	1_1	1_1	1_1	1_1	1_2	1_2	1_3	1_4	1_5	1_6	1_7	1_{1}^{2}	1^{2}_{1}	1_{1}^{2}	1_{1}^{2}	1_{1}^{2}	1_{1}^{2}	1^{3}_{1}	1^{3}_{1}	1^{3}_{1}	1_{1}^{3}	1_{1}^{3}	1_{1}^{3}
2_1	2^{1}_{1}	2^{1}_{1}	2^{1}_{1}	2^{1}_{1}	2^{1}_{1}	2^{1}_{2}	2^{1}_{2}	2^{1}_{2}	2^{1}_{2}	2^{1}_{2}	2^{1}_{2}	2^{1}_{1}	2_1	2_1	2_1	2_2	2_3	2^{3}_{1}	2^{3}_{1}	2^{3}_{1}	2^{3}_{1}	2^{3}_{1}	2^{3}_{1}
3_1	3^{1}_{1}	3^{1}_{1}	3_{1}^{1}	3^{1}_{1}	3_{1}^{1}	3^{1}_{2}	3^{1}_{2}	3^{1}_{2}	3^{1}_{2}	3^{1}_{2}	3^{1}_{2}	3_{1}^{2}		3_{1}^{2}	3_{1}^{2}	3_{1}^{2}	-	3^{1}_{1}	3_1	3_1	3_1	3_1	3_2
4_{1}^{1}	4^{1}_{1}	4^{1}_{1}	4^{1}_{1}	4^{1}_{1}	4^{1}_{1}	4^{1}_{2}	4^{1}_{2}	4^{1}_{2}	4^{1}_{2}	4^{1}_{2}	4^{1}_{2}	4_1^2	4_1^2	4_1^2	4_1^2	4_1^2	4_1^2	4^{3}_{1}	4^{3}_{1}	4_1^3	4_1^3	4_1^3	4_1^3

Figure 3: D = 6, $\rho = (5, 2, 1)$. On top: Three (D, ρ_i) -hypergraphs $H^{(1)}$, $H^{(2)}$, $H^{(3)}$. In the middle: The corresponding tableaux B_e . On the bottom: The tableaux \check{B}_e . The only differences between the middle and the bottom are highlighted in black and happen in columns that correspond to link vertices.

• If v is not a link vertex, then

the *i*-th entry of \check{B}_v is $i_{\ell(v)}$ (15.6)

the *j*-th entry
$$(j \neq i)$$
 of \check{B}_v is $j_{k(v)}^{(i)}$ (15.7)

• If $i \neq h$, then

the *i*-th entry of
$$\check{B}_{\zeta(i)}$$
 is $i^h_{k(v)}$ (15.8)

the *j*-th entry
$$(j \neq i)$$
 of $\check{B}_{\zeta^{(i)}}$ is $j_{k(v)}^i$ (15.9)

• Moreover,

the *h*-th entry of $\check{B}_{\zeta^{(h)}}$ is h_1 (15.10)

for $j \neq h, j \in I$, the *j*-th entry of $\check{B}_{\zeta^{(h)}}$ is j_1 (15.11)

for
$$j \neq h, j \notin I$$
, the *j*-th entry of $\check{B}_{\zeta^{(h)}}$ is j_1^h (15.12)

An example is provided in Figure 3. We quickly observe the following.

15.13 Claim. For $i \in I$ and a block edge e in $H^{(i)}$ we have that

- if $\zeta^{(i)} \notin e$, then $B_e = \check{B}_e$,
- if $\zeta^{(i)} \in e, i \neq h$, then B_e and \check{B}_e differ only in a single entry: The *i*-th entry of the column $\check{B}_{\zeta^{(i)}}$ is i_1^h instead of i_1 .

Proof. This follows from (15.8) and using (15.2) and (15.3).

15.14 Claim. In each row j of leftpart(T) there are only entries j_{ℓ} for some ℓ , or j_k^i for some k, i.

Proof. Recall (15.1). The claim now follows from combining (15.6), (15.7), (15.8), (15.9), (15.10), (15.11), and (15.12). \Box

15.15 Claim. If $i \notin I$, then no symbol i_{ℓ} appears in leftpart(T) for any ℓ . For a fixed $i \in I$, the symbol i_{ℓ} appears in leftpart(T) iff there is a vertex v in $H^{(i)}$ with $\ell(v) = \ell$. Moreover, i_{ℓ} appears exactly as many times as there are vertices v in $H^{(i)}$ with $\ell(v) = \ell$.

Proof. Consider (15.1) and observe that $\mathsf{leftpart}(T)$ is obtained by a permutation of the entries of the tableau

$$\sum_{i \in I} \sum_{e \in E_{\text{Block}}^{(i)}} B_e$$

Now use (15.4) and (15.5).

16 Construction of rightpart(T) for even D

The tableau rightpart(T) is constructed in any way (for example in a greedy fashion) such that the following constraints are satisfied:

rightpart(T) has the same shape as S,

a box in S has entry *i* iff there is some ℓ for which the corresponding box in rightpart(T) has entry i_{ℓ} , (16.2)

The symbol i_{ℓ} appears in rightpart(T) and leftpart(T) together exactly D many times, (16.3)

the symbol i_{ℓ} appears in rightpart(T) iff i_{ℓ} appears in leftpart(T). (16.4)

Such a tableau might not be unique, but we only care about its existence. The existence can be shown as follows.

Let $n(i_{\ell})$ denote the number of times the symbol i_{ℓ} appears in leftpart(T). If $n(i_{\ell}) > 0$, then Claim 15.15 implies that there are $n(i_{\ell}) > 0$ many vertices v in $H^{(i)}$ with $\ell(v) = i$. By Def. 14.1((4)) we know that $n(i_{\ell}) < D$. We construct rightpart(T) by arbitrarily replacing $D - n(i_{\ell})$ many entries i in S by the symbol i_{ℓ} for each i, ℓ for which $n(i_{\ell}) > 0$. Claim 16.5 below shows that this procedure replaces exactly all entries of S (recall that i appears in S exactly $D\varrho_i$ many times). It is clear that this construction satisfies (16.1), (16.2) and (16.3). Since $0 < n(i_{\ell}) < D$ iff $0 < D - n(i_{\ell}) < D$, we conclude (16.4).

16.5 Claim.

$$\forall i \in I: \quad \sum_{\ell \text{ with } n(i_{\ell}) > 0} (D - n(i_{\ell})) = D\varrho_i.$$

Proof. Since $H^{(i)}$ satisfies Def. 14.1(5) we have that

$$|E_{\text{Name}}^{(i)}| - |E_{\text{Block}}^{(i)}| = \varrho_i$$

and hence

$$D|E_{\text{Name}}^{(i)}| - D|E_{\text{Block}}^{(i)}| = D\varrho_i.$$
(*)

Moreover Def. 14.1(3) states that block edges form a set partition of V and each block edge has size D. Together with the fact that the name edges form a set partition of V (Def. 14.1(4)) we see that

(16.1)

$$\begin{split} D|E_{\text{Block}}^{(i)}| &= \sum_{e \in E_{\text{Name}}^{(i)}} \mathsf{size}(e). \text{ Together with } (*) \text{ we obtain } D|E_{\text{Name}}^{(i)}| - \sum_{e \in E_{\text{Name}}^{(i)}} \mathsf{size}(e) = D\varrho_i \text{ and hence} \\ &\sum (D - \mathsf{size}(e)) = D\varrho_i. \end{split}$$

$$\sum_{\in E_{\text{Name}}^{(i)}} (D - \mathsf{size}(e)) = D\varrho_i$$

Since for each vertex v in a name edge e the value $\ell(v)$ is the same, we write $\ell(e) := \ell(v)$. From Claim 15.15 we know that for all $e \in E_{\text{Name}}^{(i)}$ we have $n(i_{\ell(e)}) = \text{size}(e)$. Therefore

$$\sum_{e \in E_{\text{Name}}^{(i)}} (D - n(i_{\ell(e)})) = D\varrho_i$$

All numbers $\ell(e)$, $e \in E_{\text{Name}}^{(i)}$, are distinct by definition. Hence all symbols $i_{\ell(e)}$ are distinct. All $i_{\ell(e)}$ satisfy $n(i_{\ell(e)}) > 0$ by Claim 15.15. Moreover, for each ℓ with $n(i_{\ell}) > 0$ there exists some e with $\ell(e) = \ell$ also by Claim 15.15. Therefore we can rewrite the sum as

$$\sum_{\ell \text{ with } n(i_{\ell}) > 0} (D - n(i_{\ell})) = D\varrho_i,$$

which concludes the proof.

An example of the whole construction can be seen in Figure 4. We draw some quick corollaries.

e

16.6 Claim. If $i \notin I$, then the symbol i_{ℓ} does not appear in T for any ℓ . For a fixed $i \in I$, the symbol i_{ℓ} appears in T iff i_{ℓ} appears in leftpart(T) iff i_{ℓ} appears in rightpart(T) iff there is a vertex v in $H^{(i)}$ with $\ell(v) = \ell$.

Proof. We combine Claim 15.15 and (16.4).

16.7 Claim. If a symbol
$$j_k^i$$
 appears in T, then it appears exactly D many times in T.

Proof. By (16.2) the symbols j_k^i only appear in leftpart(T). Consider (15.1) and observe that leftpart(T) is obtained by a permutation of the entries of the tableau

$$\sum_{i \in I} \sum_{e \in E_{\text{Block}}^{(i)}} B_e.$$

Now use Def. 14.1((3)), (15.4), and (15.5).

17 Proof of the Tableau Lifting Theorem 13.1 for even D

In this section we prove the Tableau Lifting Theorem 13.1 for even D.

First we observe that the shape of T is indeed the required shape: This follows from Proposition 14.2, (15.1), and the fact that B_e and \check{B}_e have the same rectangular shape $m \times D$.

We remark that every symbol in T appears exactly D many times: For the symbols j_k^i this follows from Claim 16.7. For the symbols i_{ℓ} this follows from (16.3).

It remains to prove the parts (1), (2), and (3) of Theorem 13.1. We start with part (3), then build up insights that then eventually lead to the proof of parts (1) and (2).

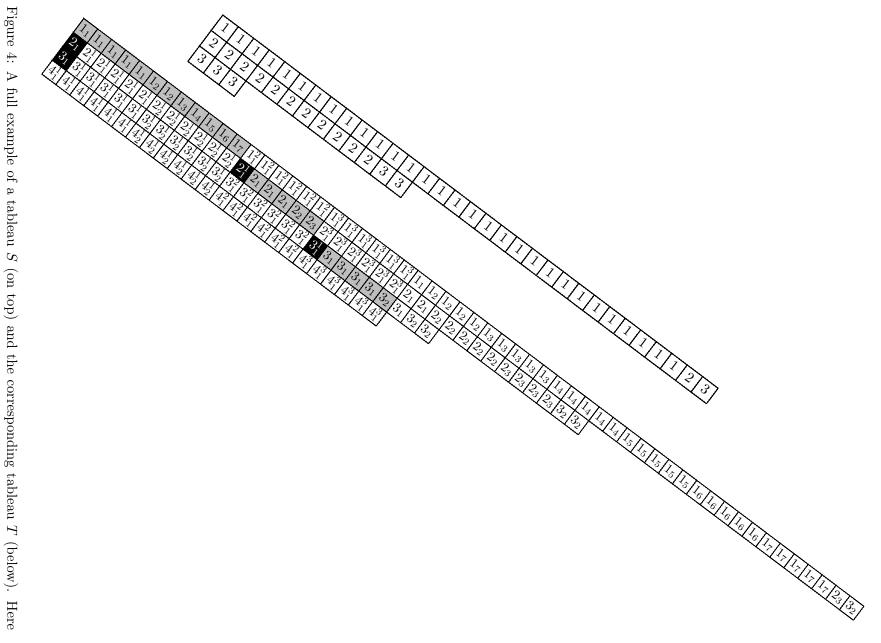


Figure 4: A full example of a tableau S (on top) and the corresponding tableau T (below). Here D = 6, m = 4.

Proof of part (3) **of Theorem 13.1**

Part (3) of Theorem 13.1 is proved as follows. We choose $\varphi(i_{\ell}) := i$ and $\varphi(j_k^i) := j$. We observe that rightpart($\varphi(T)$) = S, see (16.2). Since S is regular, rightpart($\varphi(T)$) is regular. It remains to show that leftpart($\varphi(T)$) is also regular. From Claim 15.14 we see that every column of leftpart($\varphi(T)$) contains all entries $1, \ldots, m$, sorted from top to bottom. Thus leftpart($\varphi(T)$) is regular. Since leftpart($\varphi(T)$) and rightpart($\varphi(T)$) are both regular, we conclude that $\varphi(T)$ is regular, which finishes the proof of part (3) of Theorem 13.1.

Parts (1) and (2) of Theorem 13.1: Preliminaries

In order to prove parts (1) and (2) of Theorem 13.1, we start with some preliminary observations.

17.1 Claim. If $\varphi(T)$ is regular, then for each $i \in I$ we have: For every i_{ℓ} that appears in T, $\varphi(i_{\ell})$ only depends on i and does not depend on ℓ .

Proof. By definition, for every name edge e in $H^{(i)}$ the values $\ell(v)$ coincide for all $v \in e$. This trivially implies that

for every name edge e in $H^{(i)}$: the values $\varphi(i_{\ell(v)})$ coincide for all $v \in e$. (17.2)

We claim that

for every block edge e in $H^{(i)}$: the values $\varphi(i_{\ell(v)})$ coincide for all $v \in e$. (17.3)

Proof of (17.3): Let k := k(e). According to Def. 14.1((5)) there exists a vertex $\xi^{(i)} \neq \zeta^{(i)}$ that has the same name edge and block edge as $\zeta^{(i)}$, i.e., $\ell(\zeta^{(i)}) = \ell(\xi^{(i)})$ and $k = k(\zeta^{(i)}) = k(\xi^{(i)})$. For each $v \in e, v \neq \zeta^{(i)}$ we have that the *j*-th entry $(j \neq i)$ of $\check{B}_v^{(i)}$ is j_k^i , see (15.7). Moreover, the symbol that appears as the *i*-th entry of $\check{B}_v^{(i)}$ is $i_{\ell(v)}$, see (15.6). By construction of *T*, we have that $\check{B}_v^{(i)}$ is a column in *T*. Since by assumption $\varphi(T)$ is regular, it follows that $\varphi(\check{B}_v^{(i)})$ is regular. Hence if $v \neq \zeta^{(i)}$, the $\varphi(j_k^i)$ are pairwise distinct. Thus $\varphi(i_{\ell(v)})$ equals the one element in $\{1,\ldots,m\}\setminus\{\varphi(j_k^i)\mid 1\leq j\leq m, \ j\neq i\}$. This is independent of ℓ . Hence the values $\varphi(i_{\ell(v)})$ coincide for all $v \in e, v \neq \zeta^{(i)}$. This proves (17.3) for all $v \in e, v \neq \zeta^{(i)}$. Now, if $\zeta^{(i)} \in e$, then $\xi^{(i)} \in e$, for which we have $\ell(\zeta^{(i)}) = \ell(\xi^{(i)})$, and thus clearly $\varphi(i_{\ell(\zeta^{(i)})}) = \varphi(i_{\ell(\xi^{(i)})})$. This proves the claim (17.3).

Since $H^{(i)}$ is connected (Def. 14.1(1)), we conclude with (17.2) and (17.3): The values $\varphi(i_{\ell(v)})$ coincide for all v in $H^{(i)}$. Since the symbol i_{ℓ} appears in T iff there is some vertex v in $H^{(i)}$ with $\ell(v) = \ell$ (see Claim 16.6), Claim 17.1 follows.

For $i \in I$ we define

$$\varphi^{\circ}(i) := \varphi(i_1). \tag{17.4}$$

This definition is natural, because we saw in Claim 17.1 that if $\varphi(T)$ is regular, then

$$\varphi^{\circ}(i) = \varphi(i_1) = \varphi(i_2) = \dots$$

17.5 Claim. Let $\varphi(T)$ be regular. Let $i, j \in I, i \neq j$. Then $\varphi^{\circ}(i) \neq \varphi^{\circ}(j)$.

Proof. The column $\check{B}_{\zeta^{(h)}}$ contains the symbol i_1 in row i and the symbol j_1 in row j, see (15.10) and (15.11). The fact that $\varphi(T)$ is regular implies that $\varphi(i_1) \neq \varphi(j_1)$. By (17.4) this concludes the proof.

17.6 Claim. Let $\varphi(T)$ be regular. Let $i \in I$, $i \neq h$. Then $\varphi(i_1^h) = \varphi(i_1) = \varphi^{\circ}(i)$.

Proof. The last equality is (17.4). We now prove the first equality. Let e be the block edge in $H^{(i)}$ that contains the link vertex $\zeta^{(i)}$. Then $\check{B}_e^{(i)}$ is an $m \times D$ subtableau of leftpart(T), which differs from $B_e^{(i)}$ only in a single entry in the length m column corresponding to $\zeta^{(i)}$: The *i*-th entry of the column $\check{B}_{\zeta^{(i)}}$ is i_1^h instead of i_1 , see Claim 15.13. Hence $\varphi(B_{\zeta^{(i)}})$ and $\varphi(\check{B}_{\zeta^{(i)}})$ are columns that coincide in all but at most this single box. Since $\varphi(T)$ is regular and the $\varphi(T)$ only contains entries from $\{1, \ldots, m\}$ and the columns $\varphi(B_{\zeta^{(i)}})$ and $\varphi(\check{B}_{\zeta^{(i)}})$ are of length m, we conclude that $\varphi(i_1^h) = \varphi(i_1)$.

Proof of part (1) **of Theorem 13.1**

We now prove part (1) of Theorem 13.1. The tableau rightpart(T) only contains entries i_{ℓ} and no entries j_k^i , see (16.2). As also seen in (16.2), if rightpart(T) contains an entry i_{ℓ} , then the corresponding entry of S is i. Therefore $\varphi^{\circ}(S) = \varphi(\text{rightpart}(T))$, where we lifted the map $\varphi^{\circ} : I \rightarrow$ $\{1, \ldots, m\}$ to a map with the same name that is defined on tableaux with entries from I. Claim 17.5 proves property (1) of Theorem 13.1.

Proof of part (2) **of Theorem 13.1**

The rest of this section is devoted to proving part 2 of Theorem 13.1. A rectangular tableau whose columns all coincide is called *uniform*. In the following proof we will crucially use that a uniform tableau with an even number of columns is duplex. Indeed, we prove part 2 of Theorem 13.1 by showing that if $\varphi(T)$ is regular, then for every block edge e:

- (I) $\varphi(\check{B}_e)$ is uniform if e does not contain any link vertex $\zeta^{(i)}$,
- (II) $\varphi(\check{B}_e)$ is uniform if $\zeta^{(i)} \in e$ for $i \neq h$, and
- (III) $\varphi(\check{B}_e)$ is uniform if $\zeta^{(h)} \in e$.

It is clear that these three properties cover all cases and hence $\varphi(T)$ is uniform by construction (15.1). This implies part 2 of Theorem 13.1.

We start with proving (I).

17.7 Claim. Let $\varphi(T)$ be regular. Given a block edge e in $H^{(i)}$. For all $v \in e, v \neq \zeta^{(i)}$, we have that the *i*-th entry of $\varphi(\check{B}_v)$ is $\varphi^{\circ}(i)$.

Proof. Combine (15.6) and Claim 17.1.

17.8 Claim. Let $\varphi(T)$ be regular. Given a block edge e in $H^{(i)}$. For all $j \neq i$ we have that the set

$$\{j\text{-th entry of }\varphi(\check{B}_v) \mid v \in e, v \neq \zeta^{(i)}\}$$

consists of the single element $\varphi(j_{k(e)}^i)$.

Proof. This follows from (15.7).

Combining Claim 17.7 and Claim 17.8 we see that (I) is true.

We now prove (II). Let $i \neq h$ and let e be the block edge in $H^{(i)}$ that contains $\zeta^{(i)}$. Note that k(e) = 1.

17.9 Claim. Let $\varphi(T)$ be regular. Then $\varphi(\check{B}_{\zeta(i)})$ coincides with $\varphi(\check{B}_v), v \in e, i \neq h$.

Proof. We compare the columns entrywise. Note that $k(v) = k(\zeta^{(i)}) = 1$. We make a case distinction.

Case 1: Let $j \neq i$. The *j*-th entry of \check{B}_v is j_1^i , see (15.7). The *j*-th entry of $\check{B}_{\zeta^{(i)}}$ is j_1^i , see (15.12). Hence the *j*-th entry of $\varphi(\check{B}_v)$ equals the *j*-th entry of $\varphi(\check{B}_{\zeta^{(i)}})$.

Case 2: The *i*-th entry of \check{B}_v is i_1 , see (15.6). The *i*-th entry of $\check{B}_{\zeta^{(i)}}$ is i_1^h , see (15.8). Hence Claim 17.6 implies that the *i*-th entry of $\varphi(\check{B}_v)$ equals the *i*-th entry of $\varphi(\check{B}_{\zeta^{(i)}})$.

It follows from Claim 17.9 that all columns in $\varphi(\check{B}_e)$ coincide, i.e., $\varphi(\check{B}_e)$ is uniform. Thus (II) is proved.

It remains to show (III), i.e., that $\varphi(\check{B}_e)$ is uniform if $\zeta^{(h)} \in e$.

17.10 Claim. Let $\varphi(T)$ be regular and $\zeta^{(h)}$ the link vertex in the block edge e. For all $v \in e$, $v \neq \zeta^{(h)}$, we have that the h-th entry of $\varphi(\check{B}_v)$ is $\varphi^{\circ}(h)$.

Proof. This is a direct implication of Claim 17.7.

17.11 Claim. Let $\varphi(T)$ be regular and $\zeta^{(h)} \in e$. Then $\varphi(\check{B}_{\zeta^{(h)}})$ coincides with $\varphi(\check{B}_v), v \in e$.

Proof. We compare the columns entrywise, considering three cases.

Case 1: We compare the *h*-th entry: According to Claim 17.10, the *h*-th entry of $\varphi(B_v)$ is $\varphi^{\circ}(h)$. According to (15.10) the *h*-th entry of $\check{B}_{\zeta^{(h)}}$ is h_1 , so the *h*-th entry of $\varphi(\check{B}_{\zeta^{(h)}})$ is $\varphi(h_1) = \varphi^{\circ}(h)$, see (17.4).

Case 2: We compare the *j*-th entry, $j \neq h$, in the case $j \notin I$: According to (15.7), the *j*-th entry of \check{B}_v is j_1^h . The *j*-th entry of $\check{B}_{\zeta^{(h)}}$ is also j_1^h , see (15.12). Therefore the *j*-th entry of $\varphi(\check{B}_v)$ equals the *j*-th entry of $\varphi(\check{B}_{\zeta^{(h)}})$.

Case 3: We compare the *j*-th entry, $j \neq h$, in the case $j \in I$: According to (15.7), the *j*-th entry of \check{B}_v is j_1^h . The *j*-th entry of $\check{B}_{\zeta^{(h)}}$ is j_1 , see (15.11). Claim 21.8 shows that the *j*-th entry of $\varphi(\check{B}_v)$ equals the *j*-th entry of $\varphi(\check{B}_{\zeta^{(h)}})$.

It follows from Claim 17.11 that all columns in $\varphi(\check{B}_e)$ coincide, i.e., $\varphi(\check{B}_e)$ is uniform. Thus (III) is proved. This finishes the proof of part 2 of Theorem 13.1.

Theorem 13.1 is now completely proved for even D.

18 The hypergraphs $H^{(i)}$ for odd D

Let $I := \{i \mid \varrho_i \neq 0\}$. In order to construct $\mathsf{leftpart}(T)$ we will first construct a so-called (D, ϱ_i) -paired-hypergraph for each $i \in I$. The number of columns in $\mathsf{leftpart}(T)$ will be precisely the number of vertices in all these hypergraphs together. So we want the hypergraphs to be as small as possible.

We will need the basic terms from section 14. Moreover, we will need the definition of a (D, K)-hypergraph (Def. 14.1).

18.1 Definition. Let D, K be integers. A (D, K)-paired-hypergraph is defined to be a (D, K)-hypergraph H = (V, E) that satisfies the following additional property:

Each block edge $e \in E_{\text{Block}}$ is paired with another block edge $\overline{e} \in E_{\text{Block}}$ such that they are connected by a name edge, i.e., there are vertices $v \in e$ and $\overline{v} \in \overline{e}$ called *bridge vertices* and a name edge $e_{\text{Name}} \in E_{\text{Name}}$ such that $v, \overline{v} \in e_{\text{Name}}$. In other words, the (18.2) set of block edges can be written as a disjoint union of sets of cardinality two such that the elements of each of the sets are connected by a name edge.

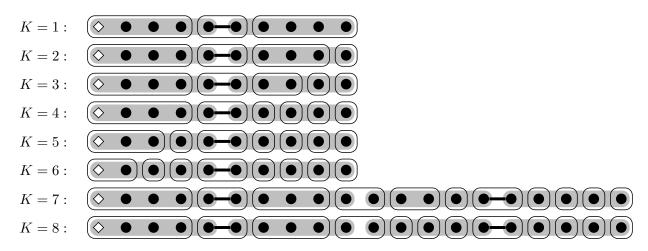


Figure 5: (5, K)-paired-hypergraphs for several values of K. Each block edges is represented as a filled gray blob. Each name edges is enclosed by a solid black curve. The link vertices are drawn as white diamonds. The bridge vertex pairs $\{v, \overline{v}\}$ for property (18.2) are illustrated via thick line segments.

Several examples for (5, K)-paired-hypergraphs are given in Figure 5. For a block edge e we write \overline{e} to denote the other block edge in its pair, and $\overline{\overline{e}} = e$.

18.3 Proposition. For odd $D \ge 3$, $K \ne 0$, there exists a (D, K)-paired-hypergraph that has exactly $2\lceil \frac{K}{2(D-2)} \rceil$ many block edges.

Proof. Let $n := 2\lceil \frac{K}{2(D-2)} \rceil$. We start the construction by considering n many disjoint block edges with D many vertices each. We arrange the vertices in a linear fashion as in Figure 5. The leftmost vertex is the link vertex. We now place these vertices in K + n many name edges as follows. As in Figure 5, the rightmost vertex of every block edge but the last shall be placed in a size 2 name edge with the leftmost vertex of the next block edge. The resulting hypergraph is connected. The rightmost vertex of every odd block edge and the leftmost vertex of every even block edge are bridge vertices. At this point we have nD - 2(n-1) = n(D-2) + 2 vertices that are not in name edges yet; and we have K + n - (n-1) = K + 1 name edges left to put vertices in. Since

$$(n(D-2)+2) - (K+1) = \left(2\left\lceil\frac{K}{2(D-2)}\right\rceil(D-2) + 2\right) - (K+1) \ge (K+2) - (K+1) = 1 > 0,$$

we can position the name edges so that the link vertex is in a name edge of size at least 2. Moreover, we position that name edge of size at least 2 in such a way that the link vertex has a vertex that not only lies in the same name edge, but also in the same block edge. \Box

19 Construction of leftpart(T) for odd D

For each $i \in I$ let $H^{(i)}$ be the (D, ϱ_i) -paired-hypergraph from Proposition 18.3. We write $E_{\text{Block}}^{(i)}$ to denote its set of block edges and $E_{\text{Name}}^{(i)}$ to denote its set of name edges.

In this section, for every $i \in I$ and every $e \in E_{\text{Block}}^{(i)}$ we construct an $m \times D$ block tableau \check{B}_e such that leftpart(T) is constructed as the concatenation

$$\mathsf{leftpart}(T) := \sum_{i \in I} \sum_{e \in E_{\mathrm{Block}}^{(i)}} \check{B}_e.$$
(19.1)

Notice that since every block edge has size D (see Def. 14.1((3))), this implies that the number of columns in leftpart(T) is equal to the sum of numbers of vertices in the hypergraphs $H^{(i)}$, $i \in I$.

Each $m \times D$ block tableau \check{B}_e is constructed in three steps: First we construct an $m \times D$ block tableau B_e , then we modify its entries to \check{B}_e , and the we make final adjustments to the entries to obtain \check{B}_e .

Let $\zeta^{(i)}$ denote the link vertex in $H^{(i)}$. We attach some additional data to each $H^{(i)}$ as follows. We put a linear order on the set of name edges $E_{name}^{(i)}$ and for each vertex v in $H^{(i)}$ we define $\ell(v)$ to be the index of its corresponding name edge. Here $\ell(v) = 1$ if v lies in the first name edge, $\ell(v) = 2$ for the next name edge, and so on. We ensure that

$$\ell(\zeta^{(i)}) = 1. \tag{19.2}$$

In the same way, we put a linear order on the set of block edge pairs; for each block edge e we write k(e) for the index of its corresponding block edge pair and for each vertex v in $H^{(i)}$ we define k(v) to be the index of its corresponding block edge pair. We ensure that

$$k(\zeta^{(i)}) = 1. \tag{19.3}$$

Moreover, for every vertex v in any $H^{(i)}$ we define i(v) := i.

In the following, for each vertex v we define an $m \times 1$ rectangular tableau (i.e., a column of length m) called B_v . Concatenating them results in B_e : $B_e := \sum_{v \in e} B_v$. The order of columns does not matter. Analogously, later we define $\dot{B}_e := \sum_{v \in e} \dot{B}_v$ and $\ddot{B}_e := \sum_{v \in e} B_v$.

Starting with B

For each block edge pair we choose one block edge to be the *barred* block edge, and the other one to be the *unbarred* block edge.

Let e be in the k-th pair of block edges in $H^{(i)}$ and let $v \in e$. The column B_v is defined by the following properties.

the *i*-th entry of
$$B_v$$
 is $i_{\ell(v)}$ (19.4)

(19.6)

the *j*-th entry
$$(j \neq i)$$
 of B_v is

$$\begin{cases} j_k^i & \text{if } e \text{ is unbarred} \\ j_{\overline{k}}^i & \text{if } e \text{ is barred} \end{cases}$$
(19.5)

An example is given in Figure 6.

From B to \hat{B}

Fix $i \in I$. To go from B_e to \dot{B}_e we switch some entries j_k^i to j_k^i and vice versa. We do this by considering the concatenation $B_e + B_{\overline{e}}$ and permuting some entries within the rows of this $m \times (2D)$ block tableau to obtain $\dot{B}_e + \dot{B}_{\overline{e}}$. For each $1 \leq k \leq |E_{\text{Block}}^{(i)}|$, let $\{e, \overline{e}\}$ denote the k-th pair of block edges in $H^{(i)}$ and choose a set of m-1 many distinct cardinality D subsets $\mathsf{barred}(i, j, k)$ of the vertex set $e \cup \overline{e}$ such that

one of the two bridge vertices is contained in all the m-1 many sets $\mathsf{barred}(i, j, k), 1 \le j \le m$, $j \ne i$, and the other bridge vertex is contained in *none* of those sets.

We define

$$\mathsf{k}\text{-}\mathsf{bar}(i,j,v) := \begin{cases} \overline{k(v)} & \text{if } v \in \mathsf{barred}(i,j,k(v)) \\ k(v) & \text{otherwise} \end{cases}$$

\diamond				•	-•)				
$1\frac{2}{1}$	$1\frac{2}{1}$	$1\frac{2}{1}$	$1\frac{2}{1}$	$1\frac{2}{1}$	1_{1}^{2}	1_{1}^{2}	1_{1}^{2}	1_{1}^{2}	1_{1}^{2}
2_1	2_{1}	2_{1}	2_{1}	2_{2}	2_2	2_3	2_3	2_3	2_4
$3\frac{2}{1}$	$3\frac{2}{1}$	$3\frac{2}{1}$	$3\frac{2}{1}$	$3\frac{2}{1}$	3_1^2	3_1^2	3_1^2	3_1^2	3_{1}^{2}
$4\frac{2}{1}$	$4\frac{2}{1}$	$4\frac{2}{1}$	$4\frac{2}{1}$	$4\frac{2}{1}$	4^{2}_{1}	4^{2}_{1}	4^{2}_{1}	4^{2}_{1}	4^{2}_{1}

Figure 6: Here i = 2 and k = 1. A (5, 2)-paired-hypergraph $H^{(2)}$ and the corresponding concatenated tableau $B_e + B_{\overline{e}}$. Vertices are drawn directly above their corresponding columns. The left block edge is considered barred, the right block edge is considered unbarred. To make the value of *i* easy to see, the second row is highlighted.

F	•	•			•		•		
$1\frac{2}{1}$									
2_1	2_{1}			2_{2}	2_{2}	2_3	2_3	2_3	2_4
$3\frac{2}{1}$	$3\frac{2}{1}$	$3\frac{2}{1}$	3_{1}^{2}			3_{1}^{2}	$3\frac{2}{1}$	3_{1}^{2}	3_{1}^{2}
$4\frac{2}{1}$	$4\frac{2}{1}$	$4\frac{2}{1}$	4^{2}_{1}	$4\frac{2}{1}$	4_1^2	4_1^2	4^{2}_{1}	4^{2}_{1}	$4\frac{2}{1}$

Figure 7: Here i = 2 and k = 1. On top: A (5, 1)-paired-hypergraph $H^{(2)}$ where the three barred sets of cardinality 5 are indicated by rectangular lines. The top rectangular line represents barred(2, 1, 1), the next line for barred(2, 2, 1) is not present (because i = 2), the next rectangular line represents barred(2, 3, 1), and the last one barred(2, 4, 1). It can be seen that the left bridge vertex is contained in all sets barred(2, ., 1), while the right bridge vertex is contained in none of those. Below: The corresponding tableau $\dot{B}_e + \dot{B}_{\overline{e}}$.

Now we define

the *i*-th entry of
$$\dot{B}_v$$
 is $i_{\ell(v)}$ (19.7)

the *j*-th entry
$$(j \neq i)$$
 of B_v is $j^i_{\mathsf{k-bar}(i,j,v)}$ (19.8)

An example is provided in Figure 7.

Since the sets $\mathsf{barred}(i, j, k)$ have cardinality D, it follows: For $i \in I$ and a block edge $e \in H^{(i)}$ we have that

 $B_e + B_{\overline{e}}$ and $\dot{B}_e + \dot{B}_{\overline{e}}$ differ only by permutations of boxes within rows, while row *i* stays the same. (19.9) For $1 \le j \le m, j \ne i$, it is now straightforward to see that

 $B_e + B_{\overline{e}}$ contains exactly the symbols j_k^i and $j_{\overline{k}}^i$, both exactly *D* many times. (19.10)

We state another insight at this point: In $H^{(i)}$, for the pair of bridge vertices v, w of the k-th block edge pair pair we have that

$$\dot{B}_v + \dot{B}_w$$
 contains all symbols from S , (19.11)

where $S := \{j_k^i \mid 1 \leq j \leq m, j \neq i\} \cup \{j_{\overline{k}}^i \mid 1 \leq j \leq m, j \neq i\}.$

Proof. This follows from combining (19.6) and (19.8): For all $j \neq i$ we have that j_k^i is the *j*-th entry of \hat{B}_v and j_k^i is the *j*-th entry of \hat{B}_w (or vice versa).

From \acute{B} to \check{B}

Most columns \dot{B}_v and \check{B}_v coincide, as we define $\check{B}_v := \dot{B}_v$ if $v \notin \{\zeta^{(i)} \mid i \in I\}$. This means that only the columns corresponding to link vertices are adjusted.

Let *h* denote the smallest number in *I*. For $i \in I$, $i \neq h$, the column $\check{B}_{\zeta^{(i)}}$ arises from $\check{B}_{\zeta^{(i)}}$ by switching the *i*-th entry with the *i*-th entry in $\check{B}_{\zeta^{(h)}}$. This means that the column $\check{B}_{\zeta^{(h)}}$ arises from $\check{B}_{\zeta^{(i)}}$ by switching the *i*-th entry with the *i*-th entry in $\check{B}_{\zeta^{(i)}}$ for all $i \in I$.

The columnwise description of B_v thus as follows.

• If v is not a link vertex, then

the *i*-th entry of \check{B}_v is $i_{\ell(v)}$ (19.12)

the *j*-th entry
$$(j \neq i)$$
 of \check{B}_v is $j_{\mathsf{k-bar}(i,j,v)}^{(i)}$ (19.13)

• If $i \neq h$, then

the *i*-th entry of $\check{B}_{\zeta^{(i)}}$ is $i^h_{\mathsf{k-bar}(h,i,\zeta^{(i)})}$ (19.14)

the *j*-th entry
$$(j \neq i)$$
 of $\mathring{B}_{\zeta^{(i)}}$ is $j^i_{\mathsf{k-bar}(i,j,\zeta^{(i)})}$ (19.15)

• Moreover,

the *h*-th entry of $\check{B}_{\zeta^{(h)}}$ is h_1 (19.16)

for $j \neq h, j \in I$, the *j*-th entry of $\dot{B}_{\zeta^{(h)}}$ is j_1 (19.17)

for
$$j \neq h, j \notin I$$
, the *j*-th entry of $\check{B}_{\zeta^{(h)}}$ is $j^{h}_{\mathsf{k}-\mathsf{bar}(h,j,\zeta^{(h)})}$ (19.18)

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1_1	1_1	1_1	1_1	1_2	1_2	1_3	1_3	1_3	1_4	1_4	1_5	1_5	1_6	1_7	1_7	1_8	1_9	1_{10}	111	$1\frac{2}{1}$	$1\frac{2}{1}$	$1\frac{2}{1}$	1_{1}^{2}	$1\frac{2}{1}$	1_{1}^{2}	$1\frac{2}{1}$	1_{1}^{2}	1_{1}^{2}	1_{1}^{2}	$1\frac{3}{1}$	$1\frac{3}{1}$	$1\frac{3}{1}$	1_{1}^{3}	$1\frac{3}{1}$	1_{1}^{3}	$1\frac{3}{1}$	1_{1}^{3}	1_{1}^{3}	1_{1}^{3}
$2\frac{1}{1}$	$2\frac{1}{1}$	$-2\frac{1}{1}$	2^{1}_{1}	$2\frac{1}{1}$	2^{1}_{1}	$2\frac{1}{1}$	2^{1}_{1}	2^{1}_{1}	2^{1}_{1}	$2\frac{1}{2}$	$2\frac{1}{2}$	$2\frac{1}{2}$	2^{1}_{2}	$2\frac{1}{2}$	2^{1}_{2}	$2\frac{1}{2}$	2^{1}_{2}	2^{1}_{2}	2^{1}_{2}	2_1	2_1	2_1	2_1	2_2	2_2	2_3	2_3	2_3	2_4	$2\frac{3}{1}$	$2\frac{3}{1}$	$2\frac{3}{1}$	2^{3}_{1}	$2\frac{3}{1}$	2_{1}^{3}	2^{3}_{1}	$2\frac{3}{1}$	2_{1}^{3}	2_{1}^{3}
$3\frac{1}{1}$	$3\frac{1}{1}$	$-3\frac{1}{1}$	3_{1}^{1}	$3\frac{1}{1}$	3_{1}^{1}	3_{1}^{1}	$3\frac{1}{1}$	3_{1}^{1}	3^{1}_{1}	$3\frac{1}{2}$	$3\frac{1}{2}$	$3\frac{1}{2}$	3^{1}_{2}	$3\frac{1}{2}$	3_{2}^{1}	3^{1}_{2}	$3\frac{1}{2}$	3^{1}_{2}	3^{1}_{2}	$3\frac{2}{1}$	$3\frac{2}{1}$	$3\frac{2}{1}$	3_{1}^{2}	$3\frac{2}{1}$	3_{1}^{2}	3_{1}^{2}	$3\frac{2}{1}$	3_{1}^{2}	3_{1}^{2}	3_1	3_1	3_1	3_1	3_2	3_2	3_3	3_3	3_3	3_3
$4\frac{1}{1}$	$4\frac{1}{1}$	$-4\frac{1}{1}$	4_{1}^{1}	$4\frac{1}{1}$	4^{1}_{1}	4_1^1	4_{1}^{1}	4^{1}_{1}	$4\frac{1}{1}$	$4\frac{1}{2}$	$4\frac{1}{2}$	$4\frac{1}{2}$	4_{2}^{1}	$4\frac{1}{2}$	4_{2}^{1}	4^{1}_{2}	4_{2}^{1}	4^{1}_{2}	$4\frac{1}{2}$	$4\frac{2}{1}$	$4\frac{2}{1}$	$4\frac{2}{1}$	4_{1}^{2}	$4\frac{2}{1}$	4_{1}^{2}	4_{1}^{2}	4_{1}^{2}	4_{1}^{2}	$4\frac{2}{1}$	$4\frac{3}{1}$	$4\frac{3}{1}$	$4\frac{3}{1}$	4_{1}^{3}	$4\frac{3}{1}$	4_1^3	4_1^3	4_1^3	4_1^3	$4\frac{3}{1}$
1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	12	12	12	12	12	12	12	12	1^{2}_{1}	12	1^{3}	1 <u>3</u>	1^{3}	13	13	13	13	13	13	13
11	11	11	11	12	12	13	13	13	14	14	15	15	16	17	17	18	19	110	111	$1\frac{2}{1}$	$1\frac{2}{1}$	$1\frac{2}{1}$	1_{1}^{2}	$1\frac{2}{1}$	1_{1}^{2}	$1\frac{2}{1}$	1_{1}^{2}	¹ 1	1_{1}^{2}	11	1	11	11	$\frac{1^{3}}{1}$	11	$1\frac{3}{1}$	11	1^{3}_{1}	11
2_1	$2\frac{1}{1}$	$-2\frac{1}{1}$	2^{1}_{1}	$2\frac{1}{1}$	2^{1}_{1}	$2\frac{1}{1}$	2^{1}_{1}	2^{1}_{1}	2^{1}_{1}	$2\frac{1}{2}$	$2\frac{1}{2}$	$2\frac{1}{2}$	2^{1}_{2}	$2\frac{1}{2}$	2^{1}_{2}	$2\frac{1}{2}$	2^{1}_{2}	2^{1}_{2}	2^{1}_{2}	$2\frac{1}{1}$	2_1	2_1	2_1	2_2	2_2	2_3	2_3	2_3	2_4	$2\frac{3}{1}$	$2\frac{3}{1}$	$2\frac{3}{1}$	2^{3}_{1}	$2\frac{3}{1}$	2^{3}_{1}	2^{3}_{1}	$2\frac{3}{1}$	2^{3}_{1}	2^{3}_{1}
3_1	$3\frac{1}{1}$	$-3\frac{1}{1}$	3^{1}_{1}	$3\frac{1}{1}$	3^{1}_{1}	3^{1}_{1}	$3\frac{1}{1}$	3^{1}_{1}	3^{1}_{1}	$3\frac{1}{2}$	$3\frac{1}{2}$	$3\frac{1}{2}$	3^{1}_{2}	$3\frac{1}{2}$	3^{1}_{2}	3^{1}_{2}	$3\frac{1}{2}$	3^{1}_{2}	3^{1}_{2}	$3\frac{2}{1}$	$3\frac{2}{1}$	$3\frac{2}{1}$	3_{1}^{2}	$3\frac{2}{1}$	3_{1}^{2}	3_{1}^{2}	$3\frac{2}{1}$	3_{1}^{2}	3_{1}^{2}	3^{1}_{1}	3_1	3_1	3_1	3_2	3_2	3_3	3_3	3_3	3_3
$4\frac{1}{1}$	$4\frac{1}{1}$	$-4\frac{1}{1}$	4_{1}^{1}	$4\frac{1}{1}$	4^{1}_{1}	4^{1}_{1}	4_1^1	4^{1}_{1}	$4\frac{1}{1}$	$4\frac{1}{2}$	$4\frac{1}{2}$	$4\frac{1}{2}$	4^{1}_{2}	$4\frac{1}{2}$	4^{1}_{2}	4^{1}_{2}	4^{1}_{2}	4^{1}_{2}	$4\frac{1}{2}$	$4\frac{2}{1}$	$4\frac{2}{1}$	$4\frac{2}{1}$	4_{1}^{2}	$4\frac{2}{1}$	4_{1}^{2}	4_{1}^{2}	4_1^2	4_{1}^{2}	$4\frac{2}{1}$	$4\frac{3}{1}$	$4\frac{3}{1}$	$4\frac{3}{1}$	4_{1}^{3}	$4\frac{3}{1}$	4_{1}^{3}	4_1^3	4_1^3	4_{1}^{3}	$4\frac{3}{1}$

Figure 8: On top: Three hypergraphs $H^{(1)}$, $H^{(2)}$, $H^{(3)}$. In the middle: The corresponding tableaux \dot{B}_e . On the bottom: The tableaux \check{B}_e . The only differences between the middle and the bottom are highlighted in black and happen in columns that correspond to link vertices.

An example is provided in Figure 8. We quickly observe the following.

19.19 Claim. For $i \in I$ and a block edge $e \in H^{(i)}$ we have that

- if $\zeta^{(i)} \notin e$, then $\dot{B}_e = \check{B}_e$,
- if $\zeta^{(i)} \in e, i \neq h$, then \dot{B}_e and \check{B}_e differ only in a single entry: The *i*-th entry of the column $\check{B}_{\zeta^{(i)}}$ is i_1^h or $i_{\overline{1}}^h$ instead of i_1 .

Proof. This follows from (19.14) and using (19.2) and (19.3).

19.20 Claim. In each row j of leftpart(T) there are only entries j_{ℓ} for some ℓ , or j_k^i or j_k^i for some k, i.

Proof. Recall (19.1). The claim now follows from combining (19.12), (19.13), (19.14), (19.15), (19.16), (19.17), and (19.18). \Box

19.21 Claim. If $i \notin I$, then no symbol i_{ℓ} appears in leftpart(T) for any ℓ . For a fixed $i \in I$, the symbol i_{ℓ} appears in leftpart(T) iff there is a vertex v in $H^{(i)}$ with $\ell(v) = \ell$. Moreover, i_{ℓ} appears exactly as many times as there are vertices v in $H^{(i)}$ with $\ell(v) = \ell$.

Proof. Consider (19.1) and observe that $\mathsf{leftpart}(T)$ is obtained by a permutation of the box entries of the tableau

$$\sum_{i \in I} \sum_{e \in E_{\text{Block}}^{(i)}} B_e$$

Now use (19.4) and (19.5).

20 Construction of rightpart(T) for odd D

The tableau rightpart(T) is constructed in any way (for example in a greedy fashion) such that the following constraints are satisfied:

rightpart(T) has the same shape as S,

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(20.1)

a box in S has entry *i* iff there is some ℓ for which the corresponding box in rightpart(T) (20.2) has entry i_{ℓ} ,

The symbol i_{ℓ} appears in rightpart(T) and leftpart(T) together exactly D many times, (20.3)

the symbol i_{ℓ} appears in rightpart(T) iff i_{ℓ} appears in leftpart(T). (20.4)

Such a tableau might not be unique, but we only care about its existence. The existence can be shown as follows.

Let $n(i_{\ell})$ denote the number of times the symbol i_{ℓ} appears in leftpart(T). If $n(i_{\ell}) > 0$, then Claim 19.21 implies that there are $n(i_{\ell}) > 0$ many vertices v in $H^{(i)}$ with $\ell(v) = i$. Using Def. 14.1((4)) we see that $n(i_{\ell}) < D$. We construct rightpart(T) by arbitrarily replacing $D - n(i_{\ell})$ many entries i in S by the symbol i_{ℓ} for each i, ℓ for which $n(i_{\ell}) > 0$. This replaces exactly all entries of S, as Claim 20.5 below shows (recall that i appears in S exactly $D\rho_i$ many times). It is clear that this construction satisfies (20.1), (20.2) and (20.3). Since $0 < n(i_{\ell}) < D$ iff $0 < D - n(i_{\ell}) < D$, we conclude (20.4).

20.5 Claim.

$$\forall i \in I: \quad \sum_{\ell \text{ with } n(i_{\ell}) > 0} (D - n(i_{\ell})) = D\varrho_i$$

Proof. Since $H^{(i)}$ satisfies Def. 14.1(5) we have that

$$|E_{\text{Name}}^{(i)}| - |E_{\text{Block}}^{(i)}| = \varrho_i$$
$$D|E_{\text{Name}}^{(i)}| - D|E_{\text{Block}}^{(i)}| = D\varrho_i.$$

and hence

Moreover Def. 14.1(3) states that block edges form a set partition of V and each block edge has size D. Together with the fact that the name edges form a set partition of V (Def. 14.1(4)) we see that
$$D|E_{\text{Block}}^{(i)}| = \sum_{e \in E_{\text{Name}}^{(i)}} \text{size}(e)$$
. Together with (*) we obtain $D|E_{\text{Name}}^{(i)}| - \sum_{e \in E_{\text{Name}}^{(i)}} \text{size}(e) = D\varrho_i$ and hence

$$\sum_{e \in E_{\text{Name}}^{(i)}} (D - \mathsf{size}(e)) = D\varrho_i.$$

Since for each vertex v in a name edge e the value $\ell(v)$ is the same, we write $\ell(e) := \ell(v)$. From Claim 19.21 we know that for all $e \in E_{\text{Name}}^{(i)}$ we have $n(i_{\ell(e)}) = \text{size}(e)$. Therefore

$$\sum_{e \in E_{\text{Name}}^{(i)}} (D - n(i_{\ell(e)})) = D\varrho_i$$

All numbers $\ell(e)$, $e \in E_{\text{Name}}^{(i)}$, are distinct by definition. Hence all symbols $i_{\ell(e)}$ are distinct. All $i_{\ell(e)}$ satisfy $n(i_{\ell(e)}) > 0$ by Claim 19.21. Moreover, for each ℓ with $n(i_{\ell}) > 0$ there exists some e with $\ell(e) = \ell$ also by Claim 19.21. Therefore we can rewrite the sum as

$$\sum_{\text{with } n(i_{\ell}) > 0} (D - n(i_{\ell})) = D\varrho_i$$

 ℓ

which concludes the proof.

(*)

An example of the whole construction can be seen in Figure 9. We draw some quick corollaries.

20.6 Claim. If $i \notin I$, then the symbol i_{ℓ} does not appear in T for any ℓ . For a fixed $i \in I$, the symbol i_{ℓ} appears in T iff i_{ℓ} appears in leftpart(T) iff i_{ℓ} appears in rightpart(T) iff there is a vertex v in $H^{(i)}$ with $\ell(v) = \ell$.

Proof. We combine Claim 19.21 and (20.4).

20.7 Claim. If a symbol j_k^i or $j_{\overline{k}}^i$ appears in T, then it appears exactly D many times in T.

Proof. By (20.2) the symbols j_k^i and j_k^i only appear in $\mathsf{leftpart}(T)$. Consider (19.1) and observe that $\mathsf{leftpart}(T)$ is obtained by a permutation of the box entries of the tableau

$$\sum_{i \in I} \sum_{e \in E_{\text{Block}}^{(i)}} B_e$$

Now use Def. 18.1((3)), (19.4), and (19.5).

21 Proof of the Tableau Lifting Theorem 13.1 for odd D

In this section we prove the Tableau Lifting Theorem 13.1 for odd D.

First we observe that the shape of T is indeed the required shape: This follows from Proposition 18.3, (19.1), and the fact that B_e and \check{B}_e have the same rectangular shape $m \times D$.

We remark that every symbol in T appears exactly D many times: For the symbols j_k^i and $j_{\overline{k}}^i$ this follows from Claim 20.7. For the symbols i_{ℓ} this follows from (20.3).

It remains to prove the parts (1), (2), and (3) of Theorem 13.1. We start with part (3), then build up insights that then eventually lead to the proof of part (1) and (2).

Proof of part (3) of Theorem 13.1

Part (3) of Theorem 13.1 is proved as follows. We choose $\varphi(i_{\ell}) := i$ and $\varphi(j_k^i) := j$ and $\varphi(j_k^i) := j$. We observe that $\operatorname{rightpart}(\varphi(T)) = S$, see (20.2). Since S is regular, $\operatorname{rightpart}(\varphi(T))$ is regular. It remains to show that $\operatorname{leftpart}(\varphi(T))$ is also regular. From Claim 19.20 we see that every column of $\operatorname{leftpart}(\varphi(T))$ contains all entries $1, \ldots, m$, sorted from top to bottom. Thus $\operatorname{leftpart}(\varphi(T))$ is regular. Since $\operatorname{leftpart}(\varphi(T))$ and $\operatorname{rightpart}(\varphi(T))$ are both regular, we conclude that $\varphi(T)$ is regular, which finishes the proof of part (3) of Theorem 13.1.

Parts (1) and (2) of Theorem 13.1: Preliminaries

In order to prove parts (1) and (2) of Theorem 13.1, we start with some preliminary observations.

21.1 Claim. If $\varphi(T)$ is regular, then for each $i \in I$, $j \neq i$, k we have: $\varphi(j_k^i) = \varphi(j_k^i)$.

Proof. Fix i, k, but do not fix j. For notational convenience we define $\dot{B} := \dot{B}_e + \dot{B}_{\overline{e}}$. Let

$$S := \{j_k^i \mid j \neq i\} \cup \{j_k^i \mid j \neq i\}$$

denote the set of symbols in \dot{B} in all rows $j \neq i$ (which is the same as the set of symbols in $B_e + B_{\overline{e}}$ in all rows $j \neq i$, see (19.9)). Note that |S| = 2(m-1) = 2m-2. Since $B_e + B_{\overline{e}}$ contains exactly

the symbols j_k^i and j_k^i , both exactly D many times (see (19.10)), the same is true for row j of \dot{B} (again, by (19.9)).

The tableau $\check{B}_e + \check{B}_{\overline{e}}$ differs from \acute{B} iff $\zeta^{(i)} \in e \cup \overline{e}$. In this case all differences are in the column that corresponds to the link vertex $\zeta^{(i)}$ (see Claim 19.19).

If $\zeta^{(i)} \in e \cup \overline{e}$, then define \dot{B}' as the $m \times (2D - 1)$ tableau that is obtained from \dot{B} by removing the column corresponding to $\zeta^{(i)}$. Note that $\zeta^{(i)}$ is not a bridge vertex. If $\zeta^{(i)} \notin e \cup \overline{e}$, then define \dot{B}' as the $m \times (2D - 1)$ tableau that is obtained from \dot{B} by removing a single arbitrary column that does not correspond to a bridge vertex.

A symbol $s \in S$ is called *abound* if it appears D many times in \dot{B}' . If $s \in S$ appears D-1 many times in \dot{B}' , then s is called *scarce*. Note that each $s \in S$ is either abound or scarce.

Note that \dot{B}' is a subtableau of T and hence $\varphi(\dot{B}')$ is regular. Let v, w denote the two bridge vertices. Since by definition they lie in the same name edge, we have $\ell(v) = \ell(w)$. Let $\ell := \ell(v)$. Note that \dot{B}_v and \dot{B}_w both have the symbol i_ℓ in row i, see (19.7). Since $\varphi(\dot{B}_v + \dot{B}_w)$ is regular and since the rows $j \neq i$ of $\dot{B}_v + \dot{B}_w$ contains all symbols from S (see (19.11)), it follows that

$$\forall s \in S: \varphi(s) \in \underbrace{\{1, \dots, m\} \setminus \{\varphi(i_\ell)\}}_{=:P}.$$
(21.2)

Note that |P| = m - 1.

Two symbols $s, t \in S$ are called *partners* if s appears exactly in those columns of \dot{B} in which t does not appear. In this case (by construction) we have $s = j_k^i$ and $t = j_k^i$, or $t = j_k^i$ and $s = j_k^i$. If two symbols s, t are not partners, then there exists a column in \dot{B} that contains both s and t, because both s and t appear in exactly D many columns in the shape $m \times (2D)$ tableau \dot{B} .

Clearly, for each partnership one partner is abound and the other is scarce.

Since $\varphi(B')$ has shape $m \times (2D - 1)$, every symbol from $\{1, \ldots, m\}$ appears exactly 2D - 1times in $\varphi(B')$. Therefore for every symbol $p \in P$ there is at most one abound symbol $s \in S$ with $\varphi(s) = p$. Since S contains m - 1 many abound symbols, from (21.2) it follows that for each $p \in P$ there is *exactly* one abound symbol $s \in S$ with $\varphi(s) = p$. Having seen this, it follows that for each $p \in P$ there is at most one scarce symbol $t \in S$ with $\varphi(s) = p$. Again, since S contains m - 1 many scarce symbols and because of (21.2), there is *exactly* one scarce symbol $t \in S$ with the property that $\varphi(t) = p$.

Given an abound symbol $s \in S$ and a symbol $t \in S$ that is not the partner of s, then in \dot{B}' there is a column that contains both s and t. Therefore $\varphi(s) \neq \varphi(t)$ if s is abound and t is not the partner of s. We conclude that for every $p \in P$ there is exactly one abound $s \in S$ and its scarce partner $t \in S$ that have $\varphi(s) = \varphi(t) = p$. This implies the claim.

21.3 Claim. If $\varphi(T)$ is regular, then for each $i \in I$ we have: For every i_{ℓ} that appears in T, $\varphi(i_{\ell})$ only depends on i and does not depend on ℓ .

Proof. By definition, for every name edge e in $H^{(i)}$ the values $\ell(v)$ coincide for all $v \in e$. This trivially implies that

for every name edge e in $H^{(i)}$: the values $\varphi(i_{\ell(v)})$ coincide for all $v \in e$. (21.4)

We claim that

for every block edge e in $H^{(i)}$: the values $\varphi(i_{\ell(v)})$ coincide for all $v \in e$. (21.5)

Proof: Let k := k(e). According to Def. 14.1((5)) there exists a vertex $\xi^{(i)} \neq \zeta^{(i)}$ that has the same name edge and block edge as $\zeta^{(i)}$, i.e., $\ell(\zeta^{(i)}) = \ell(\xi^{(i)})$ and $k = k(\zeta^{(i)}) = k(\xi^{(i)})$. For each $v \in e$,

 $v \neq \zeta^{(i)}$ we have that from each of the m-1 sets $\{j_k^i, j_k^i\}, 1 \leq j \leq m, j \neq i$, there is one symbol in the column $\check{B}_v^{(i)}$, see (19.13). Moreover, the symbol that appears as the *i*-th entry of $\check{B}_v^{(i)}$ is $i_{\ell(v)}$, see (19.12). Since $\varphi(T)$ is regular, Claim 21.1 implies that $\varphi(j_k^i) = \varphi(j_k^i)$, which we will use implicitly in the upcoming argument. For each $v \neq \zeta^{(i)}$ we have that $\check{B}_v^{(i)}$ is a column in T. In this case, since by assumption $\varphi(T)$ is regular, it follows that $\varphi(\check{B}_v^{(i)})$ is regular and hence the $\varphi(j_k^i)$ are pairwise distinct. Thus $\varphi(i_{\ell(v)})$ equals the one element in $\{1, \ldots, m\} \setminus \{\varphi(j_k^i) \mid 1 \leq j \leq m, j \neq i\}$, which is independent of ℓ . This implies that the values $\varphi(i_{\ell(v)})$ coincide for all $v \in e, v \neq \zeta^{(i)}$. This proves (21.5) for all $v \in e, v \neq \zeta^{(i)}$. Now, if $\zeta^{(i)} \in e$, then $\xi^{(i)} \in e$, for which we have $\ell(\zeta^{(i)}) = \ell(\xi^{(i)})$, and thus clearly $\varphi(i_{\ell(\zeta^{(i)})}) = \varphi(i_{\ell(\xi^{(i)})})$. This proves the claim (21.5).

Since $H^{(i)}$ is connected (Def. 14.1(1)), we conclude with (21.4) and (21.5): The values $\varphi(i_{\ell(v)})$ coincide for all v in $H^{(i)}$. Since the symbol i_{ℓ} appears in T iff there is some vertex v in $H^{(i)}$ with $\ell(v) = \ell$ (see Claim 20.6), Claim 21.3 follows.

For $i \in I$ we define

$$\varphi^{\circ}(i) := \varphi(i_1). \tag{21.6}$$

This definition is natural, because we saw in Claim 21.3 that if $\varphi(T)$ is regular, then

 $\varphi^{\circ}(i) = \varphi(i_1) = \varphi(i_2) = \dots$

21.7 Claim. Let $\varphi(T)$ be regular. Let $i, j \in I, i \neq j$. Then $\varphi^{\circ}(i) \neq \varphi^{\circ}(j)$.

Proof. The column $\check{B}_{\zeta^{(h)}}$ contains the symbol i_1 in row i and the symbol j_1 in row j, see (19.16) and (19.17). The fact that $\varphi(T)$ is regular implies that $\varphi(i_1) \neq \varphi(j_1)$. By definition (21.6), this concludes the proof.

21.8 Claim. Let $\varphi(T)$ be regular. Let $i \in I$, $i \neq h$. Then $\varphi(i_{\overline{1}}^h) = \varphi(i_1^h) = \varphi(i_1) = \varphi^{\circ}(i)$.

Proof. The first equality follows from Claim 21.1. The last equality is (21.6). We now prove the second equality. Let e be the block edge in $H^{(i)}$ that contains the link vertex $\zeta^{(i)}$. Then $\check{B}_e^{(i)}$ is an $m \times D$ subtableau of leftpart(T), which differs from $\check{B}_e^{(i)}$ only in a single entry in the length m column corresponding to $\zeta^{(i)}$: The *i*-th entry of the column $\check{B}_{\zeta^{(i)}}$ is i_1^h or $i_{\overline{1}}^h$ instead of i_1 , see Claim 19.19. Hence $\varphi(\check{B}_{\zeta^{(i)}})$ and $\varphi(\check{B}_{\zeta^{(i)}})$ are columns that coincide in all but at most this single box. Since $\varphi(T)$ is regular and the $\varphi(T)$ only contains entries from $\{1, \ldots, m\}$ and the columns $\varphi(\check{B}_{\zeta^{(i)}})$ and $\varphi(\check{B}_{\zeta^{(i)}})$ are of length m, we conclude with Claim 21.1 (i.e., $\varphi(i_1^h) = \varphi(i_{\overline{1}}^h)$) that $\varphi(i_1^h) = \varphi(i_1)$.

Proof of part (1) **of Theorem 13.1**

We now prove part (1) of Theorem 13.1. The tableau rightpart(T) only contains entries i_{ℓ} and no entries j_k^i or $j_{\overline{k}}^i$, see (20.2). As also seen in (20.2), if rightpart(T) contains an entry i_{ℓ} , then the corresponding entry of S is *i*. Therefore $\varphi^{\circ}(S) = \varphi(\text{rightpart}(T))$, where we lifted the map $\varphi^{\circ}: I \to \{1, \ldots, m\}$ to a map with the same name that is defined on tableaux with entries from *I*. Claim 21.7 proves property (1) of Theorem 13.1.

Proof of part (2) of Theorem 13.1

The rest of this section is devoted to proving part 2 of Theorem 13.1. A rectangular tableau whose columns all coincide is called *uniform*. In the following proof we will crucially use that a uniform tableau with an even number of columns is duplex. Indeed, we prove part 2 of Theorem 13.1 by showing that if $\varphi(T)$ is regular, then:

- (I) $\varphi(\check{B}_e)$ is uniform if e does not contain any link vertex $\zeta^{(i)}$,
- (II) $\varphi(\check{B}_e)$ is uniform if $\zeta^{(i)} \in e$ for $i \neq h$, and
- (III) $\varphi(\check{B}_e)$ is uniform if $\zeta^{(h)} \in e$.

It is clear that these three properties cover all cases and hence $\varphi(T)$ is uniform by construction (19.1). This implies part 2 of Theorem 13.1.

We start with proving (I).

21.9 Claim. Let $\varphi(T)$ be regular. Given a block edge e in $H^{(i)}$. For all $v \in e$, $v \neq \zeta^{(i)}$, we have that the *i*-th entry of $\varphi(\check{B}_v)$ is $\varphi^{\circ}(i)$.

Proof. Combine (19.12) and Claim 21.3.

21.10 Claim. Let $\varphi(T)$ be regular. Given a block edge e in $H^{(i)}$. For all $j \neq i$ we have that the set

$$\{j\text{-th entry of } \varphi(\check{B}_v) \mid v \in e, v \neq \zeta^{(i)}\}$$

consists of the single element $\varphi(j_{k(e)}^i)$.

Proof. Combine (19.13) and Claim 21.1.

Combining Claim 21.9 and Claim 21.10 we see that (I) is true.

We now prove (II). Let $i \neq h$ and let e be the block edge in $H^{(i)}$ that contains $\zeta^{(i)}$. Note that k(e) = 1.

21.11 Claim. Let $\varphi(T)$ be regular. Then $\varphi(\check{B}_{\zeta^{(i)}})$ coincides with $\varphi(\check{B}_v), v \in e, i \neq h$.

Proof. We compare the columns entrywise. Note that $k(v) = k(\zeta^{(i)}) = 1$. We make a case distinction.

Case 1: Let $j \neq i$. The *j*-th entry of \check{B}_v is either j_1^i or $j_{\overline{1}}^i$, see (19.13). The *j*-th entry of $\check{B}_{\zeta^{(i)}}$ is either j_1^i or $j_{\overline{1}}^i$, see (19.18). Using Claim 21.1 we see that the *j*-th entry of $\varphi(\check{B}_v)$ equals the *j*-th entry of $\varphi(\check{B}_{\zeta^{(i)}})$.

Case 2: The *i*-th entry of \check{B}_v is i_1 , see (19.12). The *i*-th entry of $\check{B}_{\zeta^{(i)}}$ is i_1^h , see (19.14). Hence Claim 21.8 implies that the *i*-th entry of $\varphi(\check{B}_v)$ equals the *i*-th entry of $\varphi(\check{B}_{\zeta^{(i)}})$.

It follows from Claim 21.11 that all columns in $\varphi(\check{B}_e)$ coincide, i.e., $\varphi(\check{B}_e)$ is uniform. Thus (II) is proved.

It remains to show (III), i.e., that $\varphi(\check{B}_e)$ is uniform if $\zeta^{(h)} \in e$.

21.12 Claim. Let $\varphi(T)$ be regular and $\zeta^{(h)}$ the link vertex in the block edge e. For all $v \in e$, $v \neq \zeta^{(h)}$, we have that the h-th entry of $\varphi(\check{B}_v)$ is $\varphi^{\circ}(h)$.

Proof. This is a direct implication of Claim 21.9.

21.13 Claim. Let $\varphi(T)$ be regular and $\zeta^{(h)} \in e$. Then $\varphi(\check{B}_{\zeta^{(h)}})$ coincides with $\varphi(\check{B}_v), v \in e$.

Proof. We compare the columns entrywise, considering three cases.

Case 1: We compare the *h*-th entry: According to Claim 21.12, the *h*-th entry of $\varphi(\check{B}_v)$ is $\varphi^{\circ}(h)$. According to (19.16) the *h*-th entry of $\check{B}_{\zeta^{(h)}}$ is h_1 , so the *h*-th entry of $\varphi(\check{B}_{\zeta^{(h)}})$ is $\varphi(h_1) = \varphi^{\circ}(h)$, see (21.6).

Case 2: We compare the *j*-th entry, $j \neq h$, in the case $j \notin I$: According to (19.13), the *j*-th entry of \check{B}_v is either j_1^h or $j_{\overline{1}}^h$. The *j*-th entry of $\check{B}_{\zeta^{(h)}}$ is j_1^h , see (19.18). Therefore Claim 21.1 shows that the *j*-th entry of $\varphi(\check{B}_v)$ equals the *j*-th entry of $\varphi(\check{B}_{\zeta^{(h)}})$.

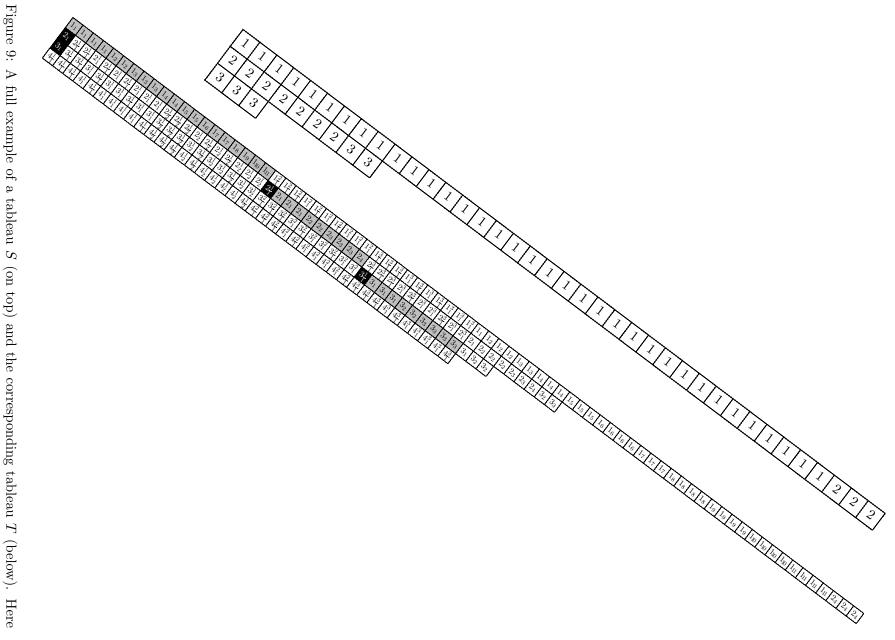
Case 3: We compare the *j*-th entry, $j \neq h$, in the case $j \in I$: According to (19.13), the *j*-th entry of \check{B}_v is either j_1^h or $j_{\overline{1}}^h$. The *j*-th entry of $\check{B}_{\zeta^{(h)}}$ is j_1 , see (19.17). Combining Claim 21.1 and Claim 21.8 shows that the *j*-th entry of $\varphi(\check{B}_v)$ equals the *j*-th entry of $\varphi(\check{B}_{\zeta^{(h)}})$.

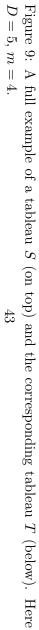
It follows from Claim 21.13 that all columns in $\varphi(\check{B}_e)$ coincide, i.e., $\varphi(\check{B}_e)$ is uniform. Thus (III) is proved. This finishes the proof of part 2 of Theorem 13.1.

Theorem 13.1 is now completely proved for odd D.

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